## Fiji Priority Rollback Protocol

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#### Situational Awareness

### Flight System















airspace data structure

Situational Awareness





#### update shared state





#### airspace data structure

Situational Awareness





Flight System

#### typically a communication protocol

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#### Situational Awareness

#### Flight System





airspace data structure



#### can we allow direct access? Situational Awareness Flight System



### Communication between mixed-criticality partitions mediate access via a lock Situational Awareness Flight System



#### Situational Awareness

#### Flight System



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#### Situational Awareness

#### Flight System





airspace data structure





#### Situational Awareness Flight System



PRP allows for priority aware, criticality aware, safe, Fill reliable, shared memory between partitions

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### PRP Locks

- Flight System's partition is guaranteed fast
- Bound on preemption based on data structure size
- Situational awareness' partition access is slightly slower but is still bounded in time



- Atomic replacement for locks
- Automatic serializability detection
- Runtime monitoring
- Aborts ability to rollback



- Atomic replicement for locks
- Automatic serializability detection
- Runtime monitoring
- Aborts ability to rollback

New programming model



- Atomic rep cement for locks
- Automatic set izability detection
- New programming model Unpredictable

- Runtime monitoring
- Aborts ability to rollback



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- Aborts ability to rollback

New programming model Unpredictable



## PRP: Two Options

- Write Buffering
  - All updates buffered : memory is always consistent
- Write Logging
  - Updates to shared memory : undo log allows reversion to consistent state of memory



synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w;
 if(foo.b+4 > foo.a)



}



synchronized(lock){
foo.a = x;
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synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w:
 if(foo.b+4 > foo.a)

}







synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w:
 if(foo.b+4 > foo.a)



Buffer



}



synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w:
 if(foo.b+4 > foo.a)







synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w:
 if(foo.b+4 > foo.a)

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synchronized(lock){
 foo.a = x;
 foo.b = y;
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synchronized(lock){
 foo.a = x;
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synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w;
 if(foo.b+4 > foo.a)







synchronized(lock) { foo.a = x; foo.b = y; foo.c = z; foo.a = w; if(foo.b+4 > foo.a)



} Commit


### Write Buffering



synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w;
 if(foo.b+4 > foo.a)

} Commit

Set lock mode to commit

foo.a , w foo.c , z foo.b , y foo.a , x

Buffer















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#### High Priority Thread

lock



WB

No priority boosting!

High Priority Thread Acquires Lock















foo.a, x



Thread

foo.a, x



















High Priority Thread Acquires Lock



### Complexity Costs - Write Buffering Low Priority Thread

- Reads: log(size of buffer) --- use RB tree
- Writes: log(size of buffer)
- Commit: size of buffer -- nested commit: n log(n)
- Acquisition: constant if thread not flushing buffer, size of buffer + context switches otherwise
- Memory: size of buffer



#### Write Logging synchronized(lock){ foo.a = x; foo.b = y;foo.c = Z; foo.a = w;if(foo.b+4 > foo.a)

foo.b

foo.c

foo.a

Main Memory



### Write Logging



synchronized(lock){
foo.a = x;
foo.b = y;
foo.c = z;
foo.a = w;
if(foo.b+4 > foo.a)

foo.a

foo.b

foo.c



Main Memory



### Write Logging



synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w;
 if(foo.b+4 > foo.a)

foo.b

foo.c

foo.a



Log

Main Memory



### Write Logging

lock set to acquired in WL mode

synchronized(lock){
 foo.a = x;
 foo.b = y;
 foo.c = z;
 foo.a = w;
 if(foo.b+4 > foo.a)

foo.a

foo.b

foo.c



Log



}



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## Write Logging synchronized(lock)

foo.a = x; foo.b = y; foo.c = z; foo.a = w; if(foo.b+4 > foo.a)

foo.b

foo.c

foo.a

## lock set to acquired in WL mode

foo.c , 3 foo.b , 2 foo.a , 1

Log

Main Memory

#### Write Logging synchronized(lock){ foo.a = x;foo.b = y;foo.c = z;foo.a = w;if(foo.b+4 > foo.a)Main foo.b foo.a foo.c Memory

## lock set to acquired in WL mode



Log



#### Write Logging synchronized(lock){ foo.a = x;foo.b = y;foo.c = Z;foo.a = w;if(foo.b+4 > foo.a)Main foo.a foo.b foo.c

lock set to acquired in WL mode

Log

Systems

foo.a, x

foo.c, 3

foo.b, 2

foo.a, 1



Memory

#### Write Logging synchronized(lock){ foo.a = x; foo.b = y;foo.c = z;foo.a = W;if(foo.b+4 > foo.a)} Main foo.b foo.a foo.c

## lock set to acquired in WL mode



Systems I

foo.a, x

foo.c, 3

foo.b , 2

foo.a, 1

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Memory

#### Write Logging synchronized(lock){ foo.a = x;foo.b = y;foo.c = Z;foo.a = w;if(foo.b+4 > foo.a)Commit

foo.a

foo.b

foo.c

#### foo.a , x foo.c , 3 foo.b , 2 foo.a , 1

lock

set to acquired

in WL mode

Log



Main Memory




















































## Complexity Costs - Write Logging Low Priority Thread

- Reads: constant
- Writes: read + log(size of write log)
- Commit: constant
- Acquisition: size of write log + context switches
- Memory: size of write log



## Questions?

