



Advancements in User Interface and Usability of KeY

Wolfram Pfeifer | September 12, 2023



KIT - The Research University in the Helmholtz Association

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KeY

Deductive verifier for (sequential) Java





KeY

Deductive verifier for (sequential) Java Java Modeling Language (JML)





KeY

Deductive verifier for (sequential) Java Java Modeling Language (JML) Modular specification/verification





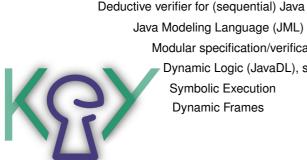
Deductive verifier for (sequential) Java Java Modeling Language (JML) Modular specification/verification Dynamic Logic (JavaDL), sequent calculus



Deductive verifier for (sequential) Java

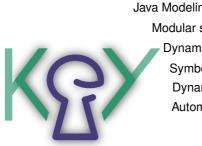
Java Modeling Language (JML) Modular specification/verification Dynamic Logic (JavaDL), sequent calculus Symbolic Execution





Java Modeling Language (JML) Modular specification/verification Dynamic Logic (JavaDL), sequent calculus Symbolic Execution **Dynamic Frames**





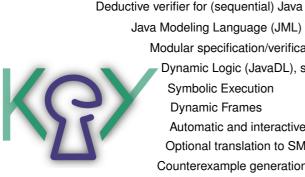
Deductive verifier for (sequential) Java Java Modeling Language (JML) Modular specification/verification Dynamic Logic (JavaDL), sequent calculus Symbolic Execution Dynamic Frames Automatic and interactive application of rules





Deductive verifier for (sequential) Java Java Modeling Language (JML) Modular specification/verification Dynamic Logic (JavaDL), sequent calculus Symbolic Execution Dynamic Frames Automatic and interactive application of rules Optional translation to SMT-LIB

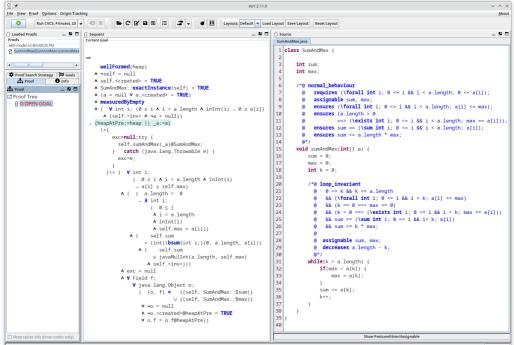


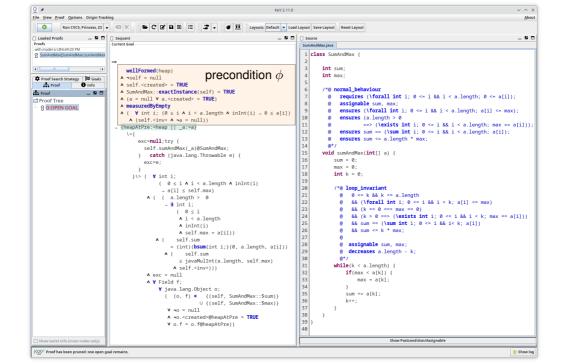


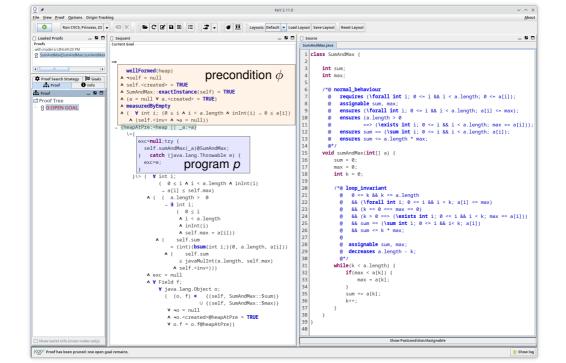
Modular specification/verification Dynamic Logic (JavaDL), sequent calculus Symbolic Execution Dynamic Frames Automatic and interactive application of rules Optional translation to SMT-LIB Counterexample generation Information flow proofs Testcase generation

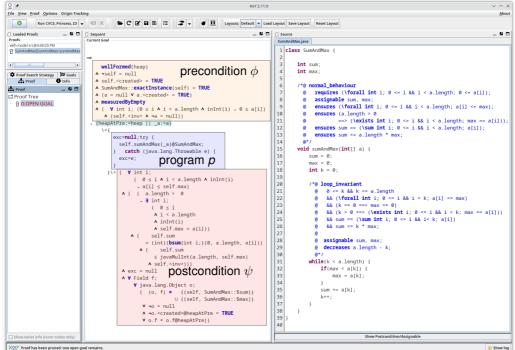
2/16 Sept. 12, 2023 W. Pfeifer: Advancements in User Interface and Usability of KeY

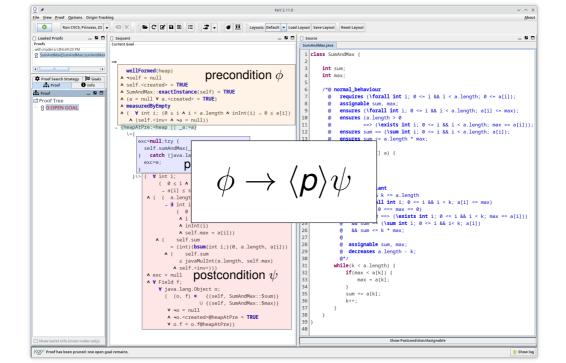
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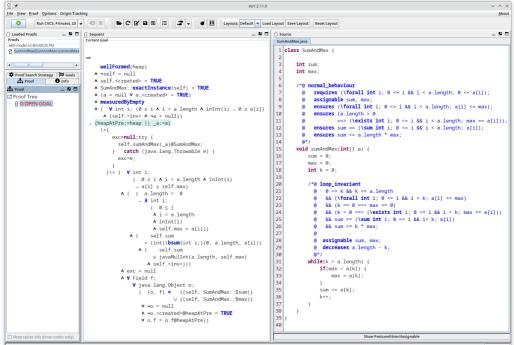


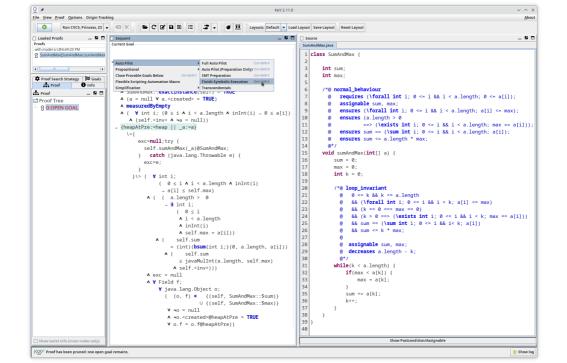


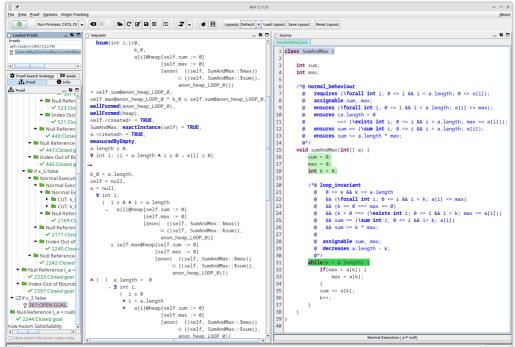


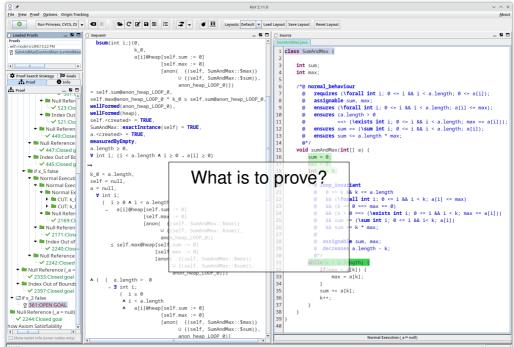


Part 1: Interaction on Source Code Level

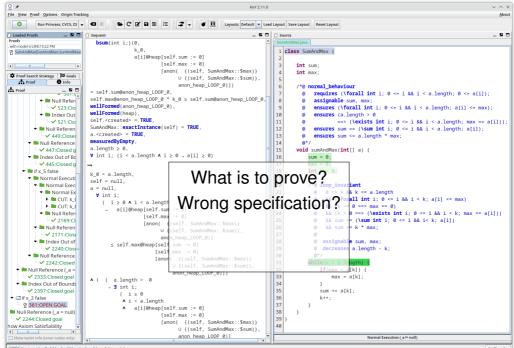


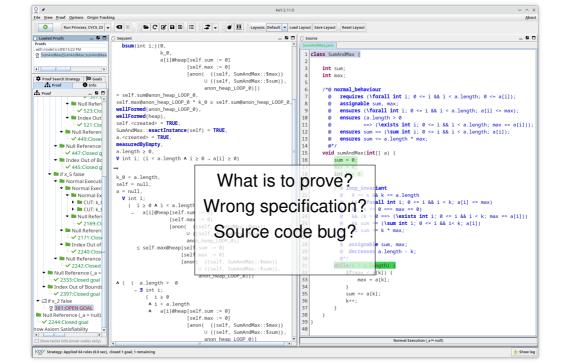


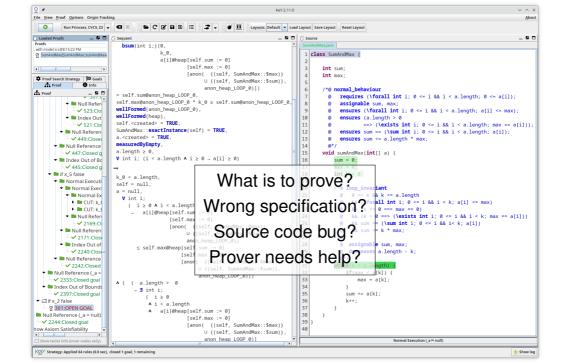


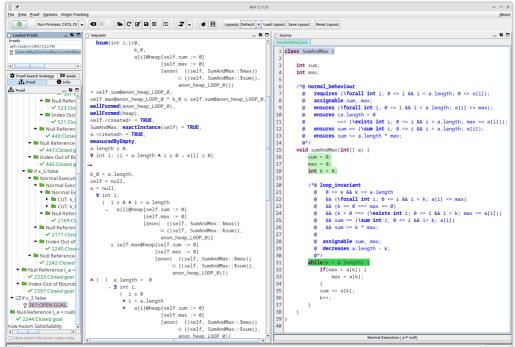


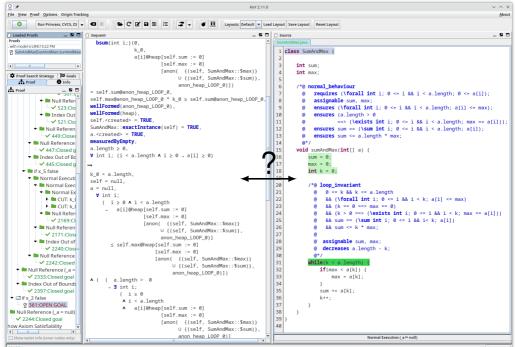
E Show log

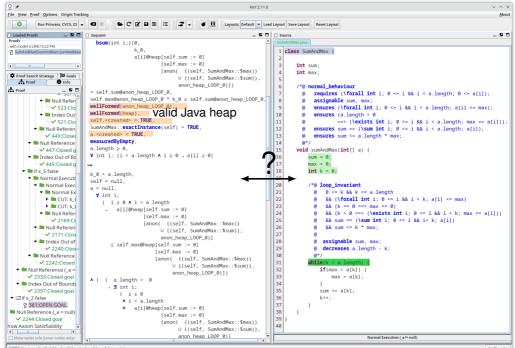


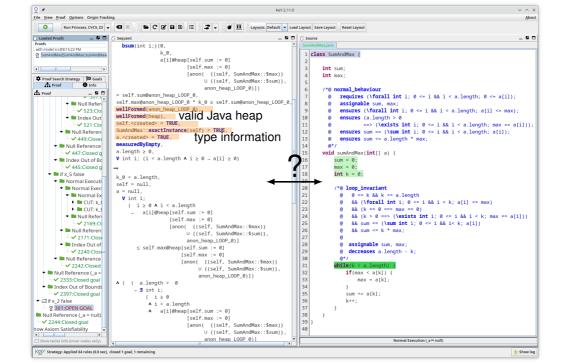


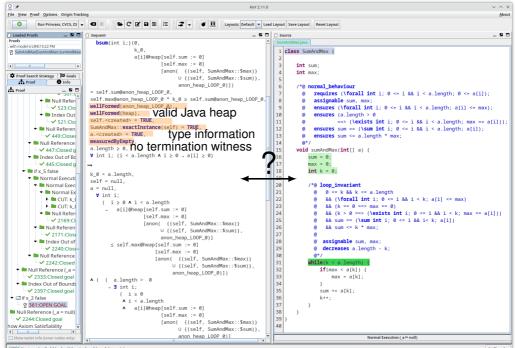


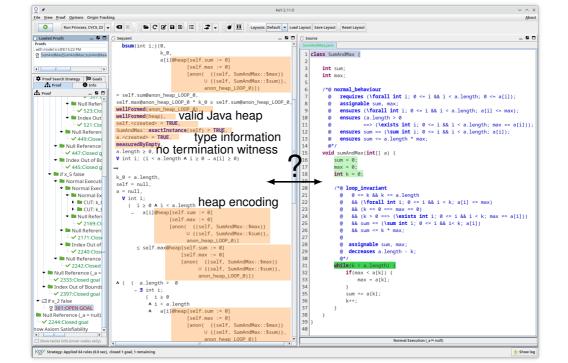












Progress so far



"Integrating Source Code, Specification and Proof State into a Single Interactive View for the Deductive Verification Tool KeY" (Master's Thesis, Mike Schwörer)

Idea: Represent a goal (sequent) of the proof as JML.

Karlsruhe Institute of Technology

Progress so far

"Integrating Source Code, Specification and Proof State into a Single Interactive View for the Deductive Verification Tool KeY" (Master's Thesis, Mike Schwörer)

Idea: Represent a goal (sequent) of the proof as JML.

- Take initial PO and assign origins/categories to the terms
- Transform correctly under rule applications
- Render the new view:

Input: Sequent with origin/category tags, Java/JML

Output: Source code with additional JML assume/assert statements placed

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Assumptions

- Symbolic execution has finished (no modalities).
- All updates are applied.
- Restrictions to allowed programs (e.g., no for loops, only return + variable, ...).

8 * KeY 2.11.0 [schwoerer/ma-ext-source-view] \vee \wedge \times File View Proof Options Origin Tracking About Run Z3. CVC5. Princess 👻 📧 💥 🖿 🕐 😰 🗄 😰 🚝 😴 🗸 💣 🎹 Layouts: Default 🗸 Load Layout: Save Layout Reset Layout 🛛 Exploration Mode 🗌 Hide justification 0 _ 8 0 C Loaded Proofs Sequent Source _ 8 0 C Loaded Proofs Cannot transform formula with modalities. - Finish symbolic execution to continue Proofs with model Part 1@3:07:05 PM CaesarChiffre.java G CaesarChiffreICaesarChiffretcalcChi 19 @ ensures (\forall int i; 0 <= i && i < valuesOutput.length; (valuesInput[i] + offset <= 'Z') ==> (valuesOutput[i] == (valuesInput[i] + offset 20 @ ensures (\forall int i; 0 <= i && i < valuesOutput.length; (valuesInput[i] + offset > 'Z') ==> (valuesOutput[i] == (valuesInput[i] + offset - 26))); 21 • Þ 22 @ ensures (\forall int i: 0 <= i && i < valuesOutput.length: 'A' <= valuesOutput[i] && valuesOutput[i] <= 'Z'); 23 ExtSourceView {{DEBUG}} 24 @ ensures \result == valuesInput.length; Goals 25 1 Info SProof Search Strategy 26 @ assignable valuesInput[*], valuesOutput[*]; Exploration Steps - Proof 27 0*/ Proof _ 8 0 int calcChiffre(int offset) { 28 29 Proof Tree 30 int loopidx = 0; ନ୍ନ 0:OPEN GOAL 31 32 convertToUpper(); 33 34 /*0 35 @ loop invariant @ <= loopidx;</pre> 36 @ loop invariant loopidx <= valuesInput.length;</pre> 37 38 @ loop_invariant (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] + offset <= 'Z') ==> (valuesOutput[i] == (valuesInput[i] + offset - 26))); 39 @ loop_invariant (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] + offset > '2') ==> (valuesOutput[i] == (valuesInput[i] + offset))); 40 41 @ loop invariant (\forall int i: 0 <= i && i < loopidx: 'A' <= valuesOutput[i] && valuesOutput[i] <= 'Z');</pre> 42 43 Ø decreases valuesInput.length - loopidx: 44 45 @ assignable valuesOutput[*]; 46 Symbolic Execution and Simplification Symbolic Execution, Simplification, and Close Provable Goals 47 while (loopidx < valuesInput.length){</pre> 48 **Close If Provable** 49 if (valuesInput[loopidx] <= 'Z' - offset) {</pre> 50 int tmp1 = valuesInput[loopidx] + offset; 51 valuesOutput[loopidx] = (char)tmp1; 52 Split and Close Provable Goals 53 } else { 54 55 int tmp2 = valuesInput[loopidx] + offset - 26; 56 valuesOutput[loopidx] = (char)tmp2; 57 58 59 60 61 loopidx++: 62 Show Postcondition/Assignable Show log 📀 KRY Replaying proof

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	36 Bloop invariant loopidx <= valuesInput.length;	
	37 8	
	38 @ loop_invariant (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] + offset <= 'Z') ==> (valuesOutput[i] == (valuesInput[i] + offset - 26));	;
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Info Proof Search Strategy	41 @ loop_invariant (\forall int i; 0 <= i && i < loopidx; 'A' <= valuesOutput[i] && valuesOutput[i] <= 'Z');	
Exploration Steps Proof	42	
	43 @ decreases valuesInput.length - loopidx;	
Proof _ 🛛 🗖	44 0	
	45 @ assignable valuesOutput[*];	
	46 e */	
lid	<pre>//@ assume \forall int i; ((i < (\old(valuesInput).length)) && (0 <= i)) ==> ('A' <= valuesInput[i] && valuesInput[i] <= 'Z');</pre>	
riant	<pre>//@ assume !(\old(valuesInput)[loopidx] <= (90 + (offset * -1)));</pre>	
n (x_arr != null)	<pre>47 while (loopidx < valuesInput.length){</pre>	
ion (x_arr_1 != null)	48	
	49 if (values:nput[loopids] <= '2' - offset) { 50 int the ! = values:nput[loopids] <= offset:	
Execution (x_arr_2 != null)		
true	51 valuesOutput[loopidx] = (char)tmp1; 52 52 52 53 53 53 53 53 53 53 53 53 53 53 53 53	
5 false ormal Execution (x arr 3 != null)	53 } else {	
Normal Execution (x_arr_5 != nul)		
Case 1	<pre>//@ assume (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] + offset <= 'Z') ==> (valuesOutput[i] == (valuesInput[i] + offset - 26)));</pre>	
+ Case 1	$//e$ assume (forall int i $0 < 1$ ba i looping, (valuesingue[] \circ offset > (2) => (valuesingue[] = (valuesingue[] \circ offset)));	
+ Case 1	//@ assume (\forall int i; 0 <= i && i < loopink; 'A' <= valueSOutput[i] && valueSOutput[i] <= '2');	1
Case 1	55 int tmp2 = valuesInput[loopidx] + offset - 26;	
→ Case 2	56 valuesOutput[loopidx] = (char)tmp2;	
R 18711:OPEN GO/	57	
Case 2	58 }	
Case 2	59	
Case 2	60	
Null Reference (x_arr_5 = null)	<pre>//@ assume offset >= 0;</pre>	
Index Out of Bounds (x_arr_5 !=)	//@ assume offset < 26;	
ull Reference (x_arr_3 = null)	<pre>//@ assume 0 <= loopidx;</pre>	
dex Out of Bounds (x_arr_3 != nul		
erence (x_arr_2 = null)	<pre>//@ assert (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] + offset > 'Z') ==> (valuesOutput[i] == (valuesInput[i] + offset)));</pre>	
ut of Bounds (x_arr_2 != null, but	62 }	
	63	
(x_arr_1 = null)	64 return valuesOutput.length;	
_arr = null)	65 }	
	66	
ertToUpper)	67 /*@ normal_behaviour	
	68 0 69 0 requires valuesInout.length > 0:	
	69 @ requires valuesInput.length > 0;	
Show taclet info (inner nodes only)	Normal Execution (x arr 5 = null)	

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		34 /*@
	(i < loopidx_0 & i >= 0	35 @ loop_invariant 0 <= loopidx;
	of bee	36 @ loop_invariant loopidx <= valuesInput.length;
	+ (self.valuesInput@heap)[i]@heapAfter_convertToUpper[anon(self.valuesOutpu	37 0
	anon_heap_LOOP)]	38 @ loop_invariant (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] +
ExtSourceView {{DEBUG}}	> 90	39 @ loop_invariant (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] +
Goals	 (self.valuesOutput@heap)[i]@heapAfter_convertToUpper[anon(self.valuesOutpu 	40 0
Info \$ Proof Search Strategy	anon_heap_LOOP)]	41 @ loop_invariant (\forall int i; 0 <= i && i < loopidx; 'A' <= valuesOutput
	= offset	42 0
Exploration Steps Proof	+ (self.valuesInput@heap)[i]@heapAfter_convertToUpper[anon(self.valuesOutp	43 @ decreases valuesInput.length - loopidx;
Proof _ 2 C	anon_heap_LOOP)	44 0
	\forall int i;	<pre>45 @ assignable valuesOutput[*];</pre>
	(i < loopidx_0 & i >= 0	46 @*/
alid	-> (self.valuesOutput@heap)[i]@heapAfter_convertToUpper[anon(self.valuesOu	<pre>//@ assume \forall int i; ((i < (\old(valuesInput).length)) && (0 <= i)) ==></pre>
ariant	anon_heap_L00	<pre>//@ assume !(\old(valuesInput)[loopidx] <= (90 + (offset * -1)));</pre>
n (x_arr != null)	>= 65	47 while (loopidx < valuesInput.length){
ution (x arr 1 != null)	& (self.valuesOutput@heap)[i]@heapAfter_convertToUpper[anon(self.valuesOu	48
and (againg the train)	anon_heap_L00	49 if (valuesInput[loopidx] <= 'Z' - offset) {
Execution (x arr 2 != null)	<= 90)	<pre>50 int tmp1 = valuesInput[loopidx] + offset;</pre>
5 true	==>	51 valuesOutput[loopidx] = (char)tmp1;
5 false	(self.valuesInput@heap)[loopidx_0]@anon_heap_convertToUpper <= 90 + offset * -1,	52
Iormal Execution (x_arr_3 l= null)	self.valuesOutput = null,	53 } else {
Normal Execution (x_arr_5 != nul	<pre>self.valuesInput = null,</pre>	54
Case 1	<pre>self.valuesOutput = self.valuesInput,</pre>	//@ assume (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] +
	self = null.	//@ assume (\forall int i; 0 <= 1 && i < loopidx; (valuesInput[i] +
Case 1	\forall int i:	<pre>//@ assume (\forall int i; 0 <= i && i < loopidx; 'A' <= valuesOutput</pre>
Case 1	(i >= 0 & i < 1 + loopidx_0	<pre>55 int tmp2 = valuesInput[loopidx] + offset - 26;</pre>
	-> offset	56 values0utput[loopidx] = (char)tmp2;
← □ Case 2	+ self.valuesInput[i]@heapAfter_convertToUpper[anon(self.valuesOutput.*,	57
Case 2	anon_heap_LOOP)]	58 }
	[self.valuesOutput[loopidx_0]	58 }
🕨 🖿 Case 2	[seriesarchaelisohia/2a]	
🖿 Case 2		60
Null Reference (x_arr_5 = null)	> 90	<pre>//@ assume offset >= 0;</pre>
Index Out of Bounds (x_arr_5 !=)	self.valuesOutput[i]@heapAfter convertToUpper[anon(self.valuesOutput.*.	//@ assume offset < 26;
ull Reference (x_arr_3 = null)	anon heap LOOP)1	<pre>//@ assume 0 <= loopidx;</pre>
dex Out of Bounds (x_arr_3 != nul	[self.valuesOutput[loopidx_0]	61 loopidx++;
erence (x_arr_2 = null)	[sell.valuesouchar[roohtav_a]	<pre>//@ assert (\forall int i; 0 <= i && i < loopidx; (valuesInput[i] + offs</pre>
ut of Bounds (x_arr_2 != null, but		62 }
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e (x_arr_1 = null)	<pre>= offset + self.valuesInput[i]@heapAfter_convertToUpper[anon(self.valuesOutput.*,</pre>	64 return valuesOutput.length;
x_arr = null)		65
	anon_heap_L00P)]	66
ertToUpper)	[self.valuesOutput[loopidx_0	67 normal_behaviour
		68
		40 menutere unlustionet length > 0.
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Part 2: Proof Slicing



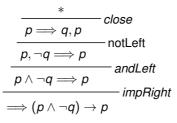
Sequent calculus: A sequent

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$$\phi_0, ..., \phi_n \implies \psi_0, ..., \psi_m$$
 is valid iff
 $\phi_0 \land ... \land \phi_n \rightarrow \psi_0 \lor ... \lor \psi_m$ is valid.

Example proof:



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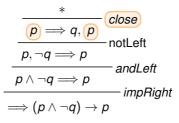
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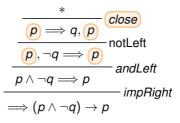
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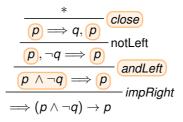
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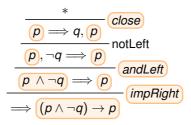
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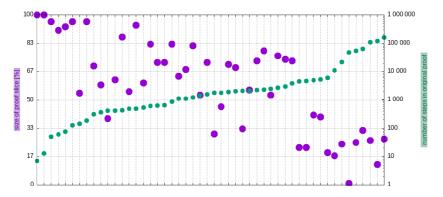
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_ Loaded Proofs Proofs	_ • •	Sequent Inner Node	_ @ U	No source loaded	_ 4
Env. with no model				No source loaded	
✓ project.key		lives(z7_0),			
· projection		killed(z7_0, agatha),			
		lives(agatha), lives(butler),			
		lives(butter),			
		z7 0 = charles,			
		<pre>\forall S z8; (llives(z8) z8 = agatha z8 = butler z8 = charles),</pre>			
Goals Proof Slicing Exploration Steps		hates(z7_0, aqatha),			
Proof D Info Proof Search	strategy	\forall S z0; (hates(z7_0, z0) !killed(charles, z0)),			
Proof	Q_80	\forall S z0; (hates(butler, z0) !killed(butler, z0)),			
	4 - 6 U	\forall S z0; (hates(aqatha, z0) Ikilled(aqatha, z0)),			
- 99:cut_direct	-	\forall S z0; (hates(charles, z0) Ikilled(charles, z0)),			
CUT: z7_0 = charles TRUE		\forall S z9; \forall S z0; (hates(z9, z0) lkilled(z9, z0)),			
 X 100:One Step Simplification: 1 rule X 122:true left 		\forall S w2; (!killed(z7_0, w2) !richer(z7_0, w2)),			
		<pre>\forall S w2; (!killed(butler, w2) !richer(butler, w2)),</pre>			
- × 123:applyEq		\forall S w1; \forall S w2; (!killed(w1, w2) !richer(w1, w2)),			
× 124:applyEq		\forall S w3; (Ihates(agatha, w3) Ihates(charles, w3)),			
× 125:applyEq		w8_1 = butler,			
→ 126:applyEq → 127:applyEq		w8 0 = butler,			
× 127:appiyeq		\forall S w4; (w4 = butler hates(agatha, w4)),			
× 128.applyEq		richer(butler, agatha),			
× 130:applyEq		hates(butler, z7_0),			
× 131:applyEq	•	<pre>\forall S w5; (hates(butler, w5) richer(w5, agatha)),</pre>			
- 132:applyEq		<pre>\forall S w6; (Ihates(agatha, w6) hates(butler, w6)),</pre>			
× 133:applyEq		\forall S w7; \exists S w8; Ihates(w7, w8)			
134:allLeft		==>			
135:replace_known_left		killed(butler, agatha),			
 136:One Step Simplification: 2 rules 		killed(butler, butler),			
137:notLeft		killed(agatha, butler),			
- 138:allLeft		killed(charles, w8_2),			
- 139:eqSymm		richer(charles, agatha),			
- 140:replace_known_right		killed(z7_0, w8_3),			
141:One Step Simplification: 2 rules		hates(z7_0, w8_3), hates(charles, w8 2),			
 142:closeFalse 		hates(charles, w8_2), hates(aqatha, butler),			
143:Closed goal		hates(agetha, butler), hates(butler, butler),			
 EUT: z7_0 = charles FALSE 		butler = agatha,			
101:One Step Simplification: 1 rule	_	killed(aqatha, aqatha)			
- 102:cut_direct		instructure adartat			
CUT: z7_0 = butler TRUE					
 X 103:One Step Simplification: 1 rule 					
- × 112:true_left					
- × 113:applyEq	-				
- X 114:applyEq					
- × 115:applyEq					
- × 116:applyEq					
 × 117:applyEq × 118:eqSymm 					
				No source load	led
- 119:applyEq	×			No source tout	

KgY Strategy: Applied 141 rules (0.3 sec), closed 3 goals, 0 remaining

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🔘 Run CVC5, Princess, Z3 👻 🗶	CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC	ayout Reset Layout 🕽 🗸 📄 Exploration Mode 🗋 Hide justification
Loaded Proofs	_ 🛛 🗖 🗋 Sequent	_ 🛛 🗖 🗋 Source 🛛 🗖
Proofs	Closed Goal	No source loaded
Env. with no model	lives(z7_0),	
✓ project.key	killed(z7_0, agatha),	
project_slice1.proof	lives(agatha),	
	lives(butler),	
	lives(charles),	
	z7 0 = charles.	
	z7_0 = aqatha z7_0 = butler true,	
Goals Proof Slicing Exploration Steps	\forall S 78: (llives(78) 78 = agatha 78 = butler 78 = charles).	
Proof O Info Proof Search	Strategy hates(charles, aqatha),	
Proof	forall S z0; (hates(z7_0, z0) !killed(z7_0, z0)),	
	\forall S z9; \forall S z0; (hates(z9, z0) !killed(z9, z0)),	
 44:replace_known_left 	\forall S w2; (lkilled(z7 0, w2) lricher(z7 0, w2)),	
 45:One Step Simplification: 2 rules 	\forall S w1; \forall S w2; (!killed(w1, w2) !richer(w1, w2)),	
- 46:notLeft	<pre>\forall S w3; (Index(agatha, w3) Index(charles, w3)),</pre>	
– 47:allLeft	w8 0 = butler,	
 48:replace_known_right 	false.	
49:One Step Simplification: 1 rule	\forall S w4; (w4 = butler hates(agatha, w4)),	
- 50:applyEq		
- 51:allLeft	hates(butler, z7_0),	
 52:replace_known_right 	<pre>\forall S w5; (hates(butler, w5) richer(w5, agatha)),</pre>	
53:One Step Simplification: 1 rule	<pre>\forall S w6; (Ihates(agatha, w6) hates(butler, w6)),</pre>	
- 54:allLeft	\forall S w7; \exists S w8; Hhates(w7, w8)	
 55:replace_known_left 	==>	
 56:One Step Simplification: 2 rules 	hates(agatha, agatha),	
- 57:cut direct	richer(z7_0, agatha),	
EUT: z7_0 = charles TRUE	hates(agatha, w8_0),	
58:applyEg	hates(butler, butler),	
- 60:allLeft	butler = agatha,	
 61:replace_known_left 	killed(agatha, agatha)	
62:One Step Simplification: 2 rules		
- 63:notLeft		
64:allLeft		
65:eqSymm		
66:replace_known_right		
67:One Step Simplification: 2 rules		
68:closeFalse		
✓ 69:Closed goal		
CUT: z7_0 = charles FALSE		
59:One Step Simplification: 1 rule		
- 70:cut_direct		
EUT: z7_0 = butler TRUE		
- 71:applyEq		
- 73:close		
✓ 74:Closed goal		
 CUT: z7 0 = butler FALSE 		
 COLL27_0 = butter FALSE 72:One Step Simplification: 1 rule 		
 72:One step simplification: 1 rule 75:applyEq 		
- 75:close		
		No source loaded
77:Closed goal		inv source loaded



Evaluation



- Often, very large parts of proofs could be removed.
- Trend: The larger the proof, the larger the percentage.
- Most of the removed steps are normalizations of formulas which are never used later on.

Part 3: Proof Caching



Motivation: Finding the correct and provable specification is often an iterative process.



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Observation: If $\Gamma \Longrightarrow \Delta$ is valid, then $\Gamma, E \Longrightarrow \Delta, Z$ is also valid (*).



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(*) Under some restrictions:

- The Java code referred to must be the same.
- Both must use the same prover settings for semantics.
- The same added rules must be present.



Motivation: Finding the correct and provable specification is often an iterative process.

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(*) Under some restrictions:

- The Java code referred to must be the same.
- Both must use the same prover settings for semantics.
- The same added rules must be present.

Ongoing work:

- Which sequents should be in the cache?
- Extend the caching beyond a single run of KeY.
- Relax the above conditions.

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] Loaded Proofs	_ S D G Sequent	_ 🛛 🗖 🗍 Source _ 🖉
Proofs with model rec(#275555 PM with model rec(#275555 PM standarded and the standard	wellFormed(heap) A self = null a self.created = TRUE A sumodbax::exactInstance(self) A (a = null) A (a = null) A (a = null) A (a = null, i < a.leng A (a = null, i < a.leng A (a = null); ty(self = null); choopAttractheap _a:=a) V < (self = null); choopAttractheap _a:=b) V < (self = null); A (a = null, i < a.leng A (a = n	7 0 requires (\forall int i; 0 <= i & i & i < a.length; 0 8 0 assignable sum assignable sum 9 ensures (\forall int i; 0 <= i & i < a.length; 0 10 ensures (\forall int i; 0 <= i & i < a.length; 1 10 ensures (\forall int i; 0 <= i & i < a.length; 1 11 0 ensures (\forall int i; 0 <= i & i < a.length; 1 12 0 ensures (a.length; 0 13 0 ensures (a.length; 0 14 0 ensures (a.length; 0 15 void sum/dbc(int i; 0 <= i & i < a.length; n 16 ensures a.length; nax; 0 17 max = 0; int k = 0; 18 int k = 0; int k = 0; 19 e 0 <= k & k & c = a.length; 10 0 e int (void and void void void and void void and void void and void void

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🕈 Edit Last Opened File 🖢	d last opened file.	Inner Node	SumAndMax.java
a) soc (Bedda) a) soc (Bedda) guiddad (Bedda) (<pre>> WellFormed(heap) A saif * null A saif * null</pre>	<pre>1 class SumArdNax { int un; int un; int un; int an; /* mormal_behaviour e require (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e exection (\formal int i; 0 <= i & i <= .length; 0 <= e & d <= & &</pre>

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Proofs	Current Goal	SumAnd Max java	
Env. with model src@7:05:05 PM	{heapAtPre:=heap		
SumAndMax(SumAndMax::sumAndMax([1)].JML normal_behavior operation	_a:=a	1 class SumAndMax {	
Env. with model src@7:10:58 PM	11 _0.00	2	
SumAndMax[SumAndMax::sumAndMax[[]]]JML normal_behavior operati	exc:=null	3 int sum;	
	II (include perturban to of fact times to of	4 int max;	
	k:=k_0	5	
	a:=k_0 * -1 + a.length)	6 /*@ normal behaviou	r
Goals Proof Slicing Exploration Steps	<pre> heap:=heap[self.sum := 0]</pre>	7 @ requires (\fo	<pre>rall int i; 0 <= 1 && 1 < a.length; 0 <=</pre>
	[self.max := 0]	8 @ assignable su	m. max:
Proof O Info Proof Search Strategy	<pre>[anon({(self, SumAndMax::\$max)}</pre>		all int i; 0 <= i && i < a.length; a[i] <
🚓 Proof 🔅 🗕 🖬 🗖	U {(self, SumAndMax::\$sum)},	10 0 ensures (a.le	
Proof Tree	anon_heap_LOOP_0)]		<pre>\exists int i; 0 <= i && i < a.length; ma</pre>
+ Invariant Initially Valid	[self.max := a[k_0]]		<pre>(\sum int i; 0 <= i && i < a.length; af</pre>
- R 42:OPEN GOAL	s_1:=self		= (vsum int i; θ <= i & i < a.iength; a[= a.length * max;
M 42:0PEN GOAL Invariant Preserved and Used	<pre> i_5:=a[k_0] + self.sum@anon_heap_LOOP_0}</pre>	13 0 ensures sun < 14 0*/	- a. rengen - max,
	(s_1 = null),		
 Normal Execution (_a != null) 	<pre>self.max@anon_heap_LOOP_0 < a[k_0],</pre>	15 void sumAndMax(int[16 sum = 0;	1 01 (
	k_0 < a.length,		
 Investigation (_a != null) 	k_0 ≥ 0,	17 max = 0;	
 	a.length ≥ k_0,	18 int k = 0;	
 Image: mail of the security of th	V int i;	19	
 Image: Second Sec	(i < k0∧i≥0	20 /*@ loop_invari	
 Image: mail of the security of th	- a[i]@heap[self.sum := 0]		k <= a.length
 Image: Image: Security of the sec			<pre>11 int i; 0 <= i && i < k; a[i] <= max)</pre>
- 🧟 352:OPEN GOAL	[self.max := 0]	23 @ && (k ==	0 ==> max == 0)
 Image: Second sec	<pre>[anon({(self, SumAndMax::\$max)}</pre>	24 @ && (k > 0	==> (\exists int i; 0 <= i && i < k; max
- 📢 337:Cached goz	<pre>U {(self, SumAndMax::\$sum)},</pre>	25 @ && sum ==	(\sum int i; 0 <= i && i < k; a[i])
 Mull Reference (_a = null) 	anon_heap_LOOP_0)]	26 @ && sum <=	k * max;
■ ■ 316:Cached goal	≤ self.max@anon_heap_LOOP_0),	27	
Index Out of Bounds (_a != null, but k (<pre>k_0 = 0 → self.max@anon_heap_LOOP_0 = 0,</pre>	28 @ assignable	sum, max:
→ ◀ 317:Cached goal	k_0 > 0		a.length - k;
 Invit Reference (s = null) 	- 3 int i;	30 @*/	
	(i < k_0	31 while(k < a.len	ath) (
 Image: Sector of the sector of	∧ i ≥ 0	32 if(max < a)	
- ← 274:Cached goal	▲ a[i]@heap[self.sum := 0]	33 max = a	
 Index Out of Bounds (_a != null, but k Out of 	[self.max := 0]	34 }	Local Action of the second s
Index Out of additions (a) = Indit, but it out of ↓ € 275:Cached goal	[anon({(self, SumAndMax::\$max)}	34 35 sum += a[k]	
if b 2 false	∪ {(self, SumAndMax::\$sum)},	36 k++;	
If D_2 faise If D_2 faise If D_2 faise If D_2 faise	anon_heap_LOOP_0)]	30 K++;	
	= self.max@anon_heap_LOOP_0),	37 }	
 Normal Execution (_a != null) 	bsum{int i;}(0,		
 Image: Image: Normal Execution (s_1 != null) 	k_8,	39 }	
କ ନୁ 267:OPEN GOAL	a[i]@heap[self.sum := 0]	40	
 Image: Image: Ima	[self.max := 0]		
4 252:Cached goal	[anon({(self, SumAndMax::\$m	11 (Y	
• I Null Reference (_a = null)	U {(self, SumAndMax::\$s		
4 234:Cached goal	anon_heap_LOOP_0)])		
 Index Out of Bounds (_a != null, but k Out 			
	<pre>= self.sum@anon_heap_LOOP_0, </pre>		
 Image: Second sec	<pre>self.max@anon_heap_LOOP_0 * k_0 ≥ self.sum@anon_heap_</pre>		n
	wellFormed(anon_heap_LOOP_0),	N	ill Reference (s_1 = null)
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Summary

We have seen:

- a novel way for interaction with KeY on the source code level.
- "Proof Slicing" to minimize sequent calculus proofs.
- "Proof Caching" to avoid reproving the same formulas.

https://github.com/KeYProject/key