

NPRG077

TinyHM: Tiny Hindley-Milner
type inference

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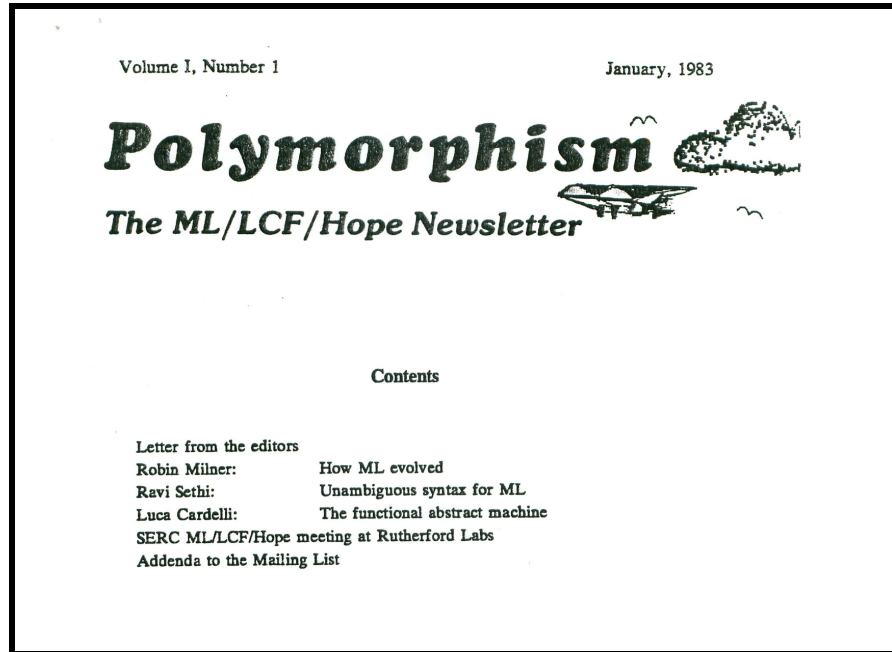


Not a programming system!?

 An important part of the ML experience
Makes ML practical and OCaml efficient

 Learn some subtle aspects of F# type inference
Some discovered late through proofs and errors

 Good example of constraint solving...
Important technique, used in Prolog & elsewhere



Origins of ML

LCF theorem prover

ML used for writing
meta-programs to
generate proofs

Types used to ensure
the validity of proofs

Hindley-Milner

A brief history of type inference

- 📖 Hindley (1969) for Combinatory Logic
- </> Milner (1978) for ML with polymorphism
- Damas (1985) with formal analysis and proofs
- 🚀 Since then - type classes, other extensions

Demo Coeffects playground

Constraint solver
code on GitHub

The screenshot shows a web browser window with the URL <https://tomasz.net/coeffects/>. The page has a navigation bar with links for Home, Theory, Papers, Short, Practice, Theory, and All. Below the navigation bar, there is a section titled "Choose a coeffect language from the dropdown and load a sample snippet to get started." This section contains a dropdown menu set to "Dataflow language (flat)" and an "Open sample" button. Below the dropdown is a code editor with the following code:

```
fun x y ->  
  let avg2 = fun y -> (y + prev y) / 2 in  
  avg2 x + prev (avg2 y)
```

Below the code editor is a "Check snippet" button. Further down, there is a text block: "In the formatted code below, you can see types of variables in a tooltip. Curried functions with multiple parameters and function defined using `let` are expanded." Below this text is another code editor with the following code:

```
fun x -> fun y ->  
  let avg2 = fun y -> (y + prev y) / 2 in  
  avg2 x + prev (avg2 y)
```

Below the code editor is another text block: "Now explore the typing derivation. Click on the judgements in the assumptions to navigate through the typing derivation. Compare flat and structural dataflow typing for the same program!" Below this text is a typing derivation diagram:

$$\frac{\frac{(\dots)}{x : \text{num}, y : \text{num} @ 1 \vdash (y + \text{prev } y) / 2 : \text{num}}}{x : \text{num} @ 1 \vdash \text{fun } y \rightarrow \dots : \text{num} \rightarrow \text{num}}}{(\dots)}$$

ML type inference

How does F# figure out the types?

Demo

Basic type inference in F#

How F# type inference works

Constraint-based

- Collect & solve constraints
- No annotations needed for ML!

```
let twice f x = f (f x)
val twice: f: ('a -> 'a) -> x: 'a -> 'a
Full name: Inference.twice
```

Let polymorphism

- Infer generic type of let-bound functions

Limitations in ML and F#

- Value restriction for generic values
- Harder to deal with .NET objects

Demo

Type inference limitations in F#

TinyHM

A bit of theory

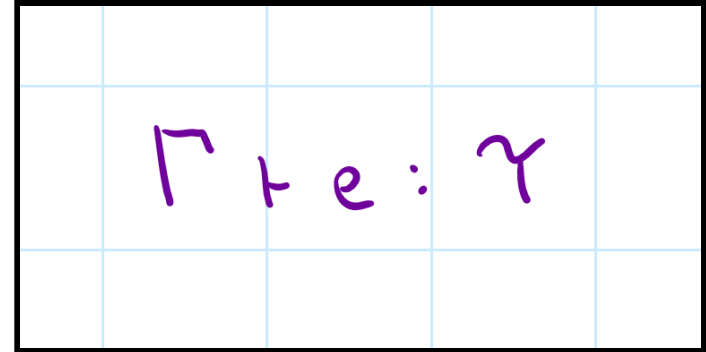
Type systems

Typing rules

Given a typing context Γ , the expression e has a type τ

The problem in general

We know some of these, want to figure out the rest



Type systems

Type checking

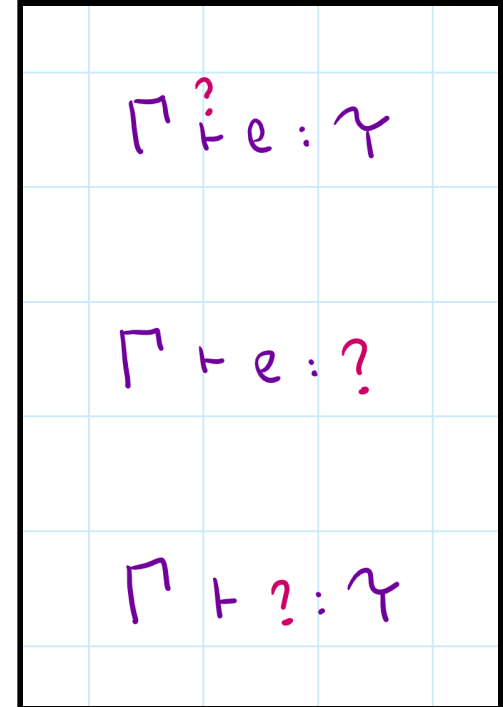
- Know it all. Check derivation exists!
- Easy for syntax-driven rules

Type inference

- Know expression. Figure out the type!
- Ideally most general (best) type

Program synthesis

- Not typical setting, but for completeness...



Type τ is the most general type
for an expression e in context Γ if

- $\Gamma \vdash e : \tau$ and
- $\forall \sigma. \Gamma \vdash e : \sigma \rightarrow \sigma <: \tau$

Principal type (most general)

Best type of
an expression

Any other type of the
expression is a special
case (subtype) of it

Type inference

- ✓ How Hindely-Milner type inference works?
Produces most general type (for ML)
- ✗ How Hindely-Milner type inference breaks?
Nominal types with members, interfaces, etc.
- ⊛ Alternative methods for type inference
Bidirectional - combines checking and inference

TinyHM

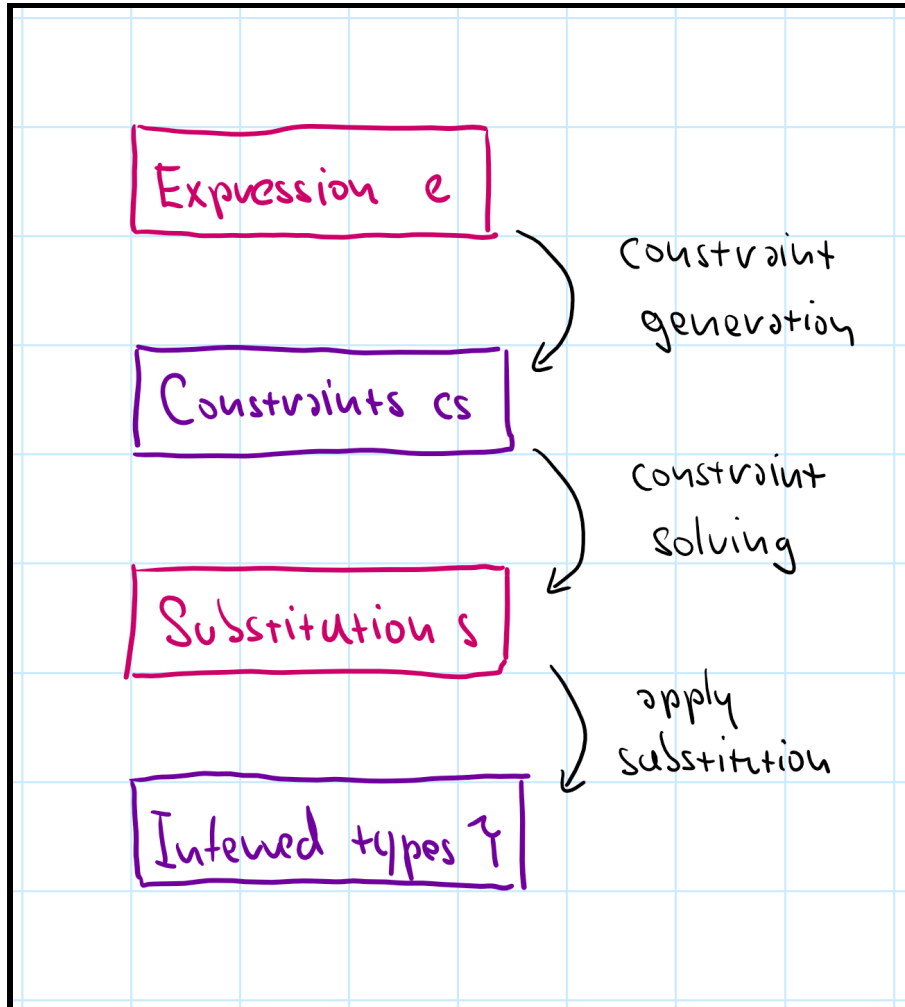
Constraint generation & solving

Two phase process

Generate constraints
Recursively over
expression

Solve constraints
Recursively over
constraint set

In the "Algorithm W",
the two are combined.
We separate them!




```
(* Basic types with
   type variables *)
type Type =
  | TyNumber
  | TyVariable of string
  | TyFunction of Type * Type
  | TyList of Type

(* Constraint specifies
   that one type should be
   unified with another *)
type Constraint =
  Type * Type
```

What is a constraint?

A pair of types that
should be unified

Easy or impossible

`int = int -> int`

`int list = int list`





Tricky with variables

`'a = int -> 'b`

`'a = 'c -> int`

TinyHM

Constraint generation

-  Generate type and constraints recursively
-  Generate new fresh type variables as needed
-  Variables with new type variables in context
-  Most checking done via constraints

Sketch

Generating constraints

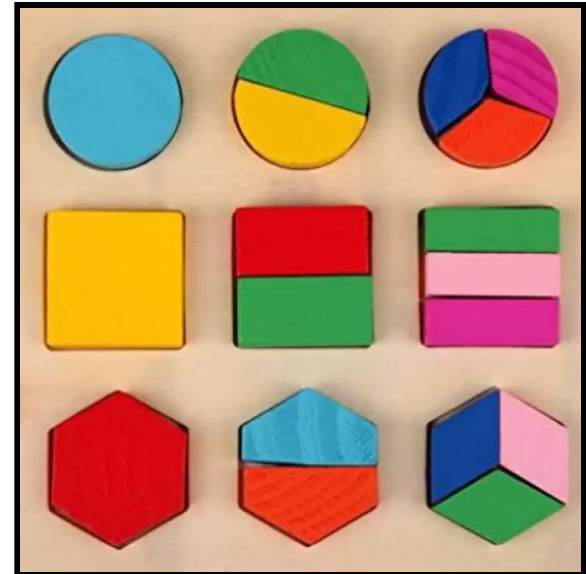
Constraint solver structure

Simplest possible example

- Peano numbers: **Zero**, **Succ(x)**
- Equality constraints with variables
- e.g. **Succ(x) = Succ(Succ(Zero))**

Creating a solver

- Discharge matching constraints
- Fail on mismatching constraints
- Generate more for matching nested
- Needs to handle substitutions...



Demo

Solving numerical constraints

Remaining work

Substitution (#1)

Replace variable in remaining constraints

Substitution (#2)

Apply substitutions to assigned type

Occurs check (#3)

Check for unsolvable constraints

```
let rec solve constraints =  
  match constraints with  
  | [] -> []  
  | (Zero, Zero)::cs -> solve cs  
  | (Succ n1, Succ n2)::cs -> solve ((n1, n2)::cs)  
  | (Zero, Succ _)::_ | (Succ _, Zero)::_ ->  
    failwith "cannot be unified"  
  | (n, Variable v)::cs | (Variable v, n)::cs ->  
    let subs = solve cs  
    (v, n)::subs
```

#1 apply all 'subs' to 'n'
#2 substitute 'n' for 'v' in remaining constraints
#3 check that 'v' does not appear in 'n'

Demo

Substitutions and occurs check

TinyHM

Inference code structure


```
(* All possible types you may
   support: type variables,
   primitives and composed *)
type Type =
  | TyVariable of string
  | TyBool
  | TyUnit
  | TyNumber
  | TyFunction of Type * Type
  | TyTuple of Type * Type
  | TyUnion of Type * Type
  | TyList of Type
  | TyForall of string * Type

(* Types of known variables *)
type TypingContext =
  Map<string, Type>
```

Types supported

Type variables

For constraint solving!

Primitive types

Match/mismatch

Composed types

Generate one or two
new constraints

Polymorphic type

Forall (bonus)

```
(* Given a list of
   constraints, produce a
   list of substitutions *)
val solve :
  list<Type * Type>
  -> list<string * Type>

(* Given a typing context
   (known variables) and
   expression, return the type
   of the expression and
   list of constraints *)
val generate :
  TypingContext
  -> Expression
  -> Type * list<Type * Type>
```

Type inference operations

Constraint solving

Takes constraints

Produces substitution

Constraint generating

Takes an expression

Produces constraints

Also check variables

Lab overview

Tiny Hindley-Milner step-by-step

TinyHM - Basic tasks

1. Complete the simple numerical constraint solver
Add the two missing substitutions to make it work!
2. Solving type constraints with numbers and Booleans
Follow the same structure, but now for type constraints...
3. Type inference for binary operators and conditionals
Add constraint generation for a subset of TinyML
4. Supporting more TinyML expressions
Add let, functions, application and occurs check
5. Adding simple data types
Constraint generation for tuples

TinyHM - Bonus & super tasks

1. Supporting more TinyML data types
Add type checking for discriminated unions
2. Type inference for lists - poor method
Add recursion & units and try this on list code!
3. Adding proper support for generic lists
New type, but without explicit type declarations
4. Inferring polymorphic code for let bindings
Implementing proper Hindley-Milner let-polymorphism
5. Exploring pathological cases
Did you know HM has DEXPTIME complexity?

Closing

Tiny Hindley-Milner type inference

Conclusions

Tiny Hindley-Milner type inference

- A remarkable quality of ML language(s)
- Cannot expect users to write types by hand!
- Nice introduction to constraint solving
- Much more can be done with this idea...

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