## Computer Architecture Digital Circuits

http://d3s.mff.cuni.cz/teaching/nswi143


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## CHARLES UNIVERSITY IN PRAGUE

faculty of mathematics and physics

## Digital computer

- Two voltage levels of interest
- High level
- Logical one (signal high/true/asserted)
- Low level
- Logical zero (signal low/false/deasserted)
- Logical values are complementary and inverse of each other
- Unlike the voltage levels representing logical ones and zeros


## Logic blocks

- Combinational
- No memory $\rightarrow$ no internal state
- Output depends only on current input
- Represents logical functions
- Sequential
- Has memory $\rightarrow$ has internal state
- Output depends on input and internal state
- Captures sequence of steps


## Logic functions and truth tables

- Logic function (also Boolean function)
- Output value is a function of input values
- $f: \mathbf{B}^{k} \rightarrow \mathbf{B}$, where $\mathbf{B}=\{0,1\}$ and $k \in \mathbf{N}$ is arity
- Truth table
- Function defined by enumerating the output for each combination of inputs (a table $2^{k}$ rows for $k$ inputs)

| Inputs |  | Output |
| :---: | :---: | :---: |
| $\boldsymbol{a}$ | $\boldsymbol{b}$ | $\boldsymbol{f}(\boldsymbol{a}, \boldsymbol{b})$ |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

## Boolean algebra

- Logic functions expressed as equations
- Variables hold values from $\mathbf{B}=\{0,1\}$
- Basic operators - primitive logic functions
- Logical inversion (NOT): $\bar{x}, \neg x,!x$
- Logical product, conjunction (AND): $x \cdot y, x \wedge y, x \& \& y$
- Logical sum, disjunction (OR): $x+y, x \vee y, x \| y$
- Additional operators (16 for 2 variables)
- NAND, NOR, XOR etc.


## Logic operators

| Inputs |  | Basic operators |  |  | Universal operators |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{a}$ | $\boldsymbol{b}$ | NOT $\boldsymbol{a}$ | $\boldsymbol{a}$ AND $\boldsymbol{b}$ | $\mathbf{a}$ OR $\boldsymbol{b}$ | $\boldsymbol{a}$ NAND $\boldsymbol{b}$ | $\boldsymbol{a}$ NOR $\boldsymbol{b}$ |
|  |  | $\urcorner$ | $\boldsymbol{\Lambda}$ | $\mathbf{y}$ | $\boldsymbol{\uparrow}$ | $\boldsymbol{\downarrow}$ |
| 0 | 0 | 1 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 | 0 | 0 |


| Inputs <br> $\boldsymbol{a}$ |  |  | $\boldsymbol{b}$ | $\boldsymbol{a}$ XOR $\boldsymbol{b}$ | $\boldsymbol{a}$ XNOR $\boldsymbol{b}$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 1 | 1 | $\ldots$ | 0 | 1 |
| 0 | 1 | 1 | 0 | 1 | 0 | $\ldots$ | 0 | 1 |
| 1 | 0 | 1 | 0 | 0 | 1 | $\ldots$ | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 | $\ldots$ | 0 | 1 |

## Boolean algebra laws

- Idempotency: $a+a=a, a \cdot a=a$
- Computativity: $a+b=b+a, a \cdot b=b \cdot a$
- Associativity: $a+(b+c)=(a+b)+c, a \cdot(b \cdot c)=(a \cdot b) \cdot c$
- Absorption: $a \cdot(a+b)=a, a+(a \cdot b)=a$
- Distributivity: $a \cdot(b+c)=(a \cdot b)+(a \cdot c), a+(b \cdot c)=(a+b) \cdot(a+c)$
- Neutrality of 0 and 1: $a+0=a, a \cdot 1=a$
- Aggressivity of 0 and 1: $a+1=1, a \cdot 0=0$
- Complementarity: $a+\neg a=1, a \cdot \neg a=0$
- Absorption of negation: $a \cdot(\neg a+b)=a \cdot b, a+(\neg a \cdot b)=a+b$
- De Morgan's laws: $\neg(a+b)=\neg a \cdot \neg b, \neg(a \cdot b)=\neg a+\neg b$
- Double negation: $\neg(\neg a)=a$


## Intermezzo: CPU logical operations (1)

- Logic functions extended to operate on (finite) sequences of bits
- Word = finite sequence of bits
- Word length = number of bits in the sequence
- Output of a logic operation a is function of input values
- $f:\left(\mathbf{B}^{n}\right)^{k} \rightarrow \mathbf{B}^{n}$, where $\mathbf{B}=\{0,1\}, k \in \mathbf{N}$ is arity and $n \in \mathbf{N}$ is word length


## Intermezzo: CPU logical operations (2)

- (Bitwise) logical product/sum/inversion
- Operators \& I, ~ etc. in C-like languages
- Primitive logic function applied to individual bits of the input words, result stored to individual bits of the output word
- Allow isolating (AND), zeroing (AND, NOR), setting (OR), inverting (XOR) selected bits, or inverting all bits (NOT), of the input word


## Intermezzo: CPU logical operations (3)

- Logical shifts (left and right)
- Operators << and >> in C-like languages
- Shifts bits in a words $i$ positions to the left or right
- "Vacated" bits are replaced with 0
- For binary natural numbers
- Shift by $\boldsymbol{i}$ bits to the left $\rightarrow$ multiplying by $\mathbf{2}^{i}$
- Shift by $\boldsymbol{i}$ bits to the right $\rightarrow$ dividing by $\mathbf{2}^{\boldsymbol{i}}$


## Logic gates (1)

- Physical implementation basic logic functions
- Basic gates: NOT, OR, AND



## Logic gates (2)

- Physical implementation of logic operators
- Inverting gates: NAND, NOR
- Less common gates: XOR



## Combinational logic circuits

## - Implementation of more complex logic functions

- Combines multiple logic operators
- Logic signals correspond to variables
- Logic gates correspond to primitive operators
- Most commonly NAND or NOR gates
- Sufficient for expressing any logic function
- Logic block
- Abstracts away from internal structure of a circuit
- Provides functional building blocks


## Logic blocks: binary (half) adder

- Adds two 1-bit numbers
- The simplest case
- Input:
- operand a
- operand b
- Output:
- sum s
- carry c
- Function:
- $s=a \cdot \neg b+\neg a \cdot b=a \operatorname{XOR} b$
- $c=a \cdot b=a$ AND $b$


| Inputs |  | Outputs |  |
| :---: | :---: | :---: | :---: |
| $\boldsymbol{a}$ | $\boldsymbol{b}$ | $\boldsymbol{c}$ |  |
| 0 | $\boldsymbol{s}$ |  |  |
| 0 | 0 | 0 |  |
|  | 1 | 0 |  |
| 1 | 0 | 0 |  |
| 1 | 1 | 1 |  |

## Logic blocks: binary (half) adder (2)



## Logic blocks: binary adder (3)

- Adding $n$-bit numbers
- Merge $\boldsymbol{n}^{1 ⁄ 2}$-adders for individual bits?

- $1 / 2$-adder cannot propagate carry from previous additions (not enough inputs)


## Logic blocks: binary adder (4)

## - Full adder

- Adds two 1-bit numbers taking into account carry from previous addition
- Input:
- operand a
- operand b
- carry $c_{0}$
- Output:
- sum s
- carry c


## Logic blocks: binary adder (5)

## - Full adder

- Adds two 1-bit numbers taking into account carry from previous addition
- Inputs: operand $\boldsymbol{a}$, operand $\boldsymbol{b}$, carry $c_{0}$

| $\boldsymbol{c}_{\mathbf{0}}$ | $\boldsymbol{a}$ | $\boldsymbol{b}$ | $\boldsymbol{c}$ |
| :---: | :---: | :---: | :---: |
| 0 | $\boldsymbol{s}$ |  |  |
| 0 | 0 | 0 | 0 |
| 0 |  |  |  |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

- Outputs: sum $s$, carry $c$

$$
\text { - } \begin{aligned}
s & =\neg c_{0} \cdot(a \cdot \neg b+\neg a \cdot b)+c_{0} \cdot(a \cdot b+\neg a \cdot \neg b) \\
s & =\ldots \\
s & =c_{0} \mathrm{XOR}(a \operatorname{XOR} b) \\
-c & =a \cdot b+c_{0} \cdot(a \cdot \neg b+\neg a \cdot b) \\
c & =(a \operatorname{AND} b) \operatorname{OR}\left(c_{0} \operatorname{AND}(a \operatorname{XOR} b)\right)
\end{aligned}
$$

## Logic blocks: binary adder (6)



## Logic blocks: binary adder (7)



## Logic blocks: binary adder (8)



## Logic block for subtraction

- Taking advantage of 2's complement
- Basic building block: adder
- Use XOR gate as a controlled inverter
- Example: 2-bit ALU supporting addition and subtraction
- Data input: operand bits $\boldsymbol{a}_{1} \boldsymbol{a}_{0}$, operand bits $\boldsymbol{b}_{1} \boldsymbol{b}_{0}$
- Control input: signal SUB to determine operation
- $S U B=0 \rightarrow$ addition
- $S U B=1 \rightarrow$ subtraction
- Output: sum/difference bits $\boldsymbol{s}_{1} \boldsymbol{s}_{0}$, carry $\boldsymbol{c}$


## 2-bit ALU for adding/subtracting



## Sequential logic

- Combinational logic + memory elements
- Memory elements keep internal state
- Inputs and the contents of memory (internal state) determines outputs and next internal state
- Synchronous vs. asynchronous sequential circuits
- Determines how and when state changes
- Need to ensure stable inputs (inputs don't change)



## Synchronous sequential circuits

- Clock signal to synchronize state changes
- Change state during one clock cycle
- Inputs of combinational logic does not change while it is being read
- Writing of values from outputs to memory elements happens with rising/falling edge of the clock signal
rising edge



## Memory elements

- Pair of inverters with feedback
- Asynchronous circuit with two stable states
- Allows "storing" 1 bit of information
- We just need to be able to control the state...
- Basic building block for memory elements



## Set-Reset (R-S) latch, NAND-based



| Inputs |  | Outputs |  |
| :---: | :---: | :---: | :---: |
| $\overline{\boldsymbol{r}}$ | $\overline{\boldsymbol{s}}$ | $\boldsymbol{q}_{\boldsymbol{n}}$ | $\overline{\boldsymbol{q}}_{\boldsymbol{n}}$ |
| $\mathbf{0}$ | 0 | $?$ | $?$ |
| $\mathbf{0}$ | 1 | 0 | 1 |
| $\mathbf{1}$ | 0 | 1 | 0 |
| $\mathbf{1}$ | 1 | $\boldsymbol{q}_{n-1}$ | $\neg \boldsymbol{q}_{n-1}$ |

## Set-Reset (R-S) latch, NOR-based



| Vstupy |  | Výstupy |  |
| :---: | :---: | :---: | :---: |
| $\boldsymbol{r}$ | $\boldsymbol{s}$ | $\boldsymbol{q}_{\boldsymbol{n}}$ | $\overline{\boldsymbol{q}}_{\boldsymbol{n}}$ |
| 0 | 0 | $q_{n-1}$ | $\neg \boldsymbol{q}_{n-1}$ |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | $?$ | $?$ |

## Other flip-flops

- Derived from R-S
- Clocked R-S latch
- Synchronous R-S latch variant
- Reacts to $r$ or $\boldsymbol{s}$ signals while the clock signal is high
- R-S master/slave (R-S flip-flop)
- Reacts to $r$ or $\boldsymbol{s}$ signals only on rising/falling edge of the clock signal



## Other flip-flops (2)

- Derived from R-S
- J-K master/slave (J-K flip-flop)
- Extends R-S (J = S, K = R), inverts state when J = K = 1
- Clocked D latch, D flip-flop
- Value determines by single input
- T flip-flop
- Allows dividing clock signal frequency



## Data register made of flip-flops



## Shift register made of flip-flops



## 32-bit sequential multiplier



## 32-bit sequential divider



