Computer Architecture Digital Circuits

http://d3s.mff.cuni.cz/teaching/nswi143



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faculty of mathematics and physics

Digital computer

Two voltage levels of interest

- High level
 - Logical one (signal high/true/asserted)
- Low level
 - Logical zero (signal low/false/deasserted)
- Logical values are complementary and inverse of each other
 - Unlike the voltage levels representing logical ones and zeros



Logic blocks

Combinational

- No memory \rightarrow no internal state
- Output depends only on current input
- Represents logical functions

Sequential

- Has memory \rightarrow has internal state
- Output depends on input and internal state
- Captures sequence of steps



Logic functions and truth tables

- Logic function (also Boolean function)
 - Output value is a function of input values
 - $f: \mathbf{B}^k \rightarrow \mathbf{B}$, where $\mathbf{B} = \{0, 1\}$ and $k \in \mathbf{N}$ is arity
 - Truth table
 - Function defined by enumerating the output for each combination of inputs (a table 2^k rows for k inputs)

Inputs		Output	
a	b	f(a, b)	
0	0	0	
0	1	1	
1	0	1	
1	1	1	



Boolean algebra

Logic functions expressed as equations

- Variables hold values from B = {0, 1}
- Basic operators primitive logic functions
 - Logical inversion (NOT): \overline{x} , $\neg x$, !x
 - Logical product, conjunction (AND): $x \cdot y$, $x \wedge y$, x & y
 - Logical sum, disjunction (OR): x + y, x V y, x || y
- Additional operators (16 for 2 variables)
 - NAND, NOR, XOR etc.



Logic operators

Inp	Inputs Basic operators		Universal operators			
a	b	NOT a	a AND b	a OR <i>b</i>	a NAND b	a NOR b
		_	٨	V	<u> </u>	\checkmark
0	0	1	0	0	1	1
0	1	1	0	1	1	0
1	0	0	0	1	1	0
1	1	0	1	1	0	0

Inp	uts	Other operators						
a	b	a XOR b	a XNOR b					
		Ð	\leftrightarrow	\rightarrow	÷	•••		Т
0	0	0	1	1	1	• • •	0	1
0	1	1	0	1	0	• • •	0	1
1	0	1	0	0	1	• • •	0	1
1	1	0	1	1	1	• • •	0	1



Boolean algebra laws

- Idempotency: a + a = a, $a \cdot a = a$
- **Computativity:** a + b = b + a, $a \cdot b = b \cdot a$
- Associativity: a + (b + c) = (a + b) + c, $a \cdot (b \cdot c) = (a \cdot b) \cdot c$
- Absorption: $a \cdot (a + b) = a, a + (a \cdot b) = a$
- **Distributivity:** $a \cdot (b + c) = (a \cdot b) + (a \cdot c), a + (b \cdot c) = (a + b) \cdot (a + c)$
- Neutrality of 0 and 1: a + 0 = a, $a \cdot 1 = a$
- Aggressivity of 0 and 1: a + 1 = 1, $a \cdot 0 = 0$
- **Complementarity:** $a + \neg a = 1$, $a \cdot \neg a = 0$
- Absorption of negation: $a \cdot (\neg a + b) = a \cdot b$, $a + (\neg a \cdot b) = a + b$
- De Morgan's laws: $\neg(a + b) = \neg a \cdot \neg b$, $\neg(a \cdot b) = \neg a + \neg b$
- Double negation: $\neg(\neg a) = a$



Intermezzo: CPU logical operations (1)

- Logic functions extended to operate on (finite) sequences of bits
 - Word = finite sequence of bits
 - Word length = number of bits in the sequence
 - Output of a logic operation a is function of input values
 - *f*: (**B**ⁿ)^k → **B**ⁿ, where **B** = {0, 1}, k∈**N** is arity and n∈**N** is word length



Intermezzo: CPU logical operations (2)

- (Bitwise) logical product/sum/inversion
 - Operators &, |, ~ etc. in C-like languages
 - Primitive logic function applied to individual bits of the input words, result stored to individual bits of the output word
 - Allow isolating (AND), zeroing (AND, NOR), setting (OR), inverting (XOR) selected bits, or inverting all bits (NOT), of the input word



Intermezzo: CPU logical operations (3)

Logical shifts (left and right)

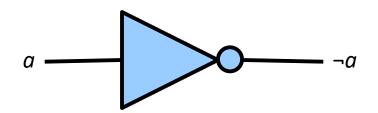
- Operators << and >> in C-like languages
- Shifts bits in a words *i* positions to the left or right
 - "Vacated" bits are replaced with 0
- For binary natural numbers
 - Shift by *i* bits to the left \rightarrow multiplying by 2^i
 - Shift by **i** bits to the right \rightarrow dividing by $\mathbf{2}^{i}$

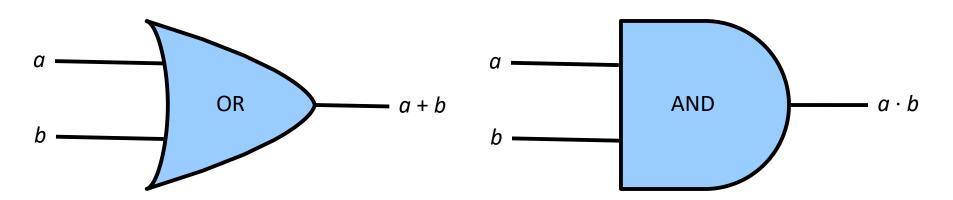


Logic gates (1)

Physical implementation basic logic functions

Basic gates: NOT, OR, AND



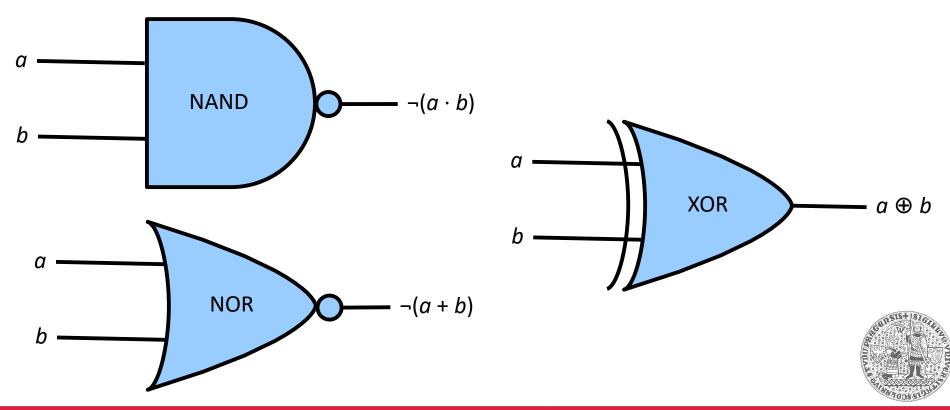




Logic gates (2)

• Physical implementation of logic operators

- Inverting gates: NAND, NOR
- Less common gates: XOR



Combinational logic circuits

Implementation of more complex logic functions

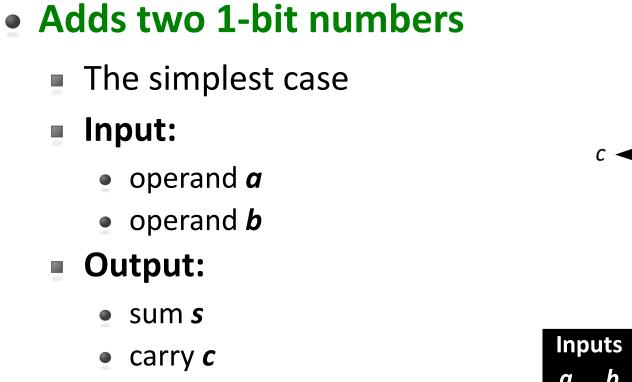
- Combines multiple logic operators
 - Logic signals correspond to variables
 - Logic gates correspond to primitive operators
- Most commonly NAND or NOR gates
 - Sufficient for expressing any logic function

Logic block

- Abstracts away from internal structure of a circuit
- Provides functional building blocks



Logic blocks: binary (half) adder



Function:

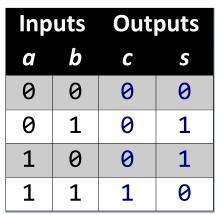
Output:

sum s

carry **c**

Input:

- $s = a \cdot \neg b + \neg a \cdot b = a \text{ XOR } b$
- $c = a \cdot b = a \text{ AND } b$



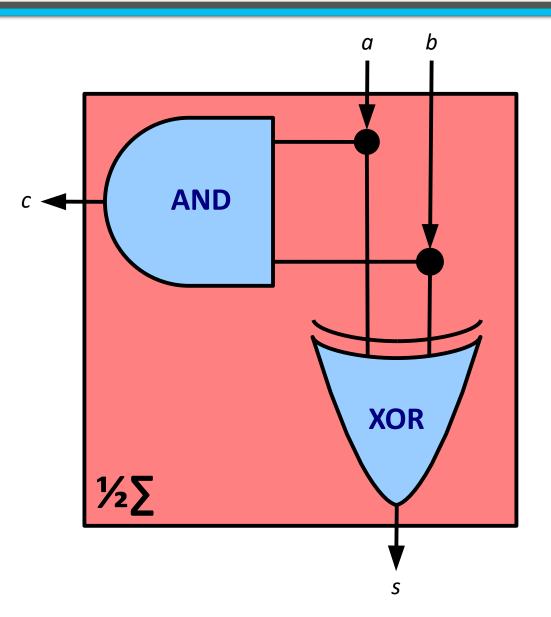


b

а

½∑

Logic blocks: binary (half) adder (2)

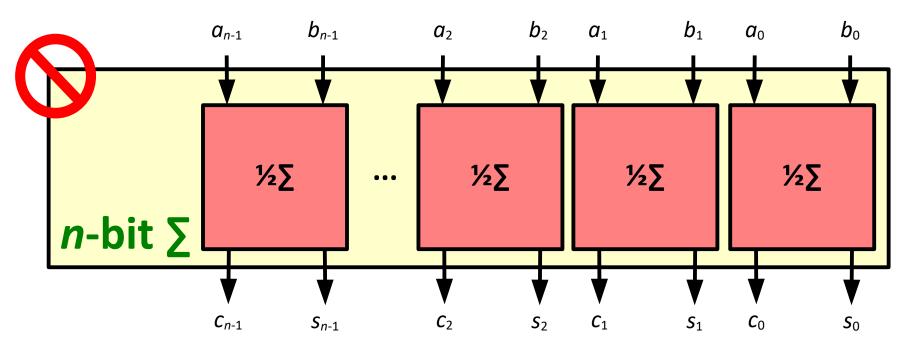




Logic blocks: binary adder (3)

• Adding *n*-bit numbers

Merge n ½-adders for individual bits?



• ½-adder cannot propagate carry from previous additions (not enough inputs)



Logic blocks: binary adder (4)

Full adder

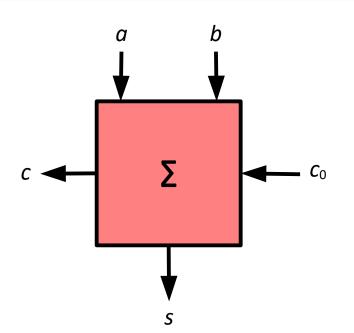
Adds two 1-bit numbers taking into account carry from previous addition

Input:

- operand *a*
- operand *b*
- e carry **c**o

Output:

- sum *s*
- carry *c*





Logic blocks: binary adder (5)

• Full adder

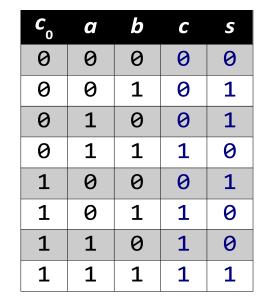
- Adds two 1-bit numbers taking into account carry from previous addition
- Inputs: operand *a*, operand *b*, carry *c*₀
- Outputs: sum s, carry c

•
$$s = \neg c_0 \cdot (a \cdot \neg b + \neg a \cdot b) + c_0 \cdot (a \cdot b + \neg a \cdot \neg b)$$

 $s = \dots$
 $s = c_0 \text{ XOR } (a \text{ XOR } b)$

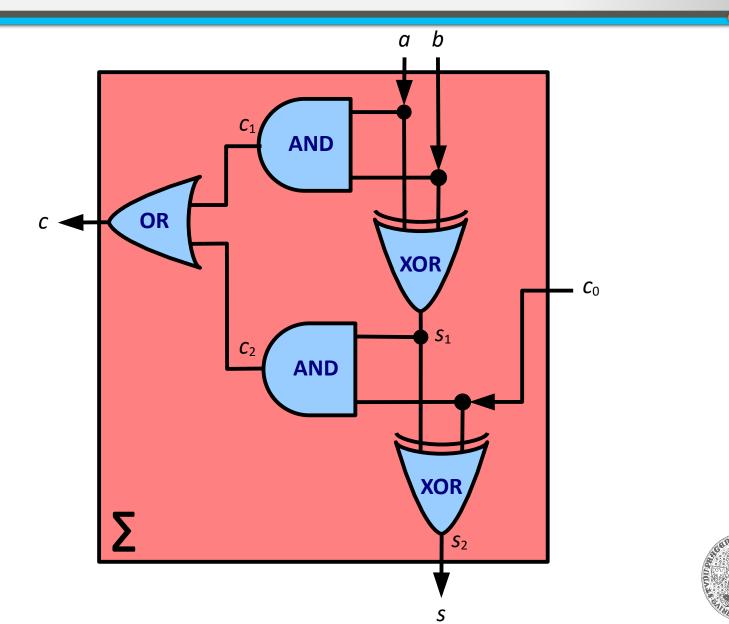
•
$$c = a \cdot b + c_0 \cdot (a \cdot \neg b + \neg a \cdot b)$$

 $c = (a \text{ AND } b) \text{ OR } (c_0 \text{ AND } (a \text{ XOR } b))$



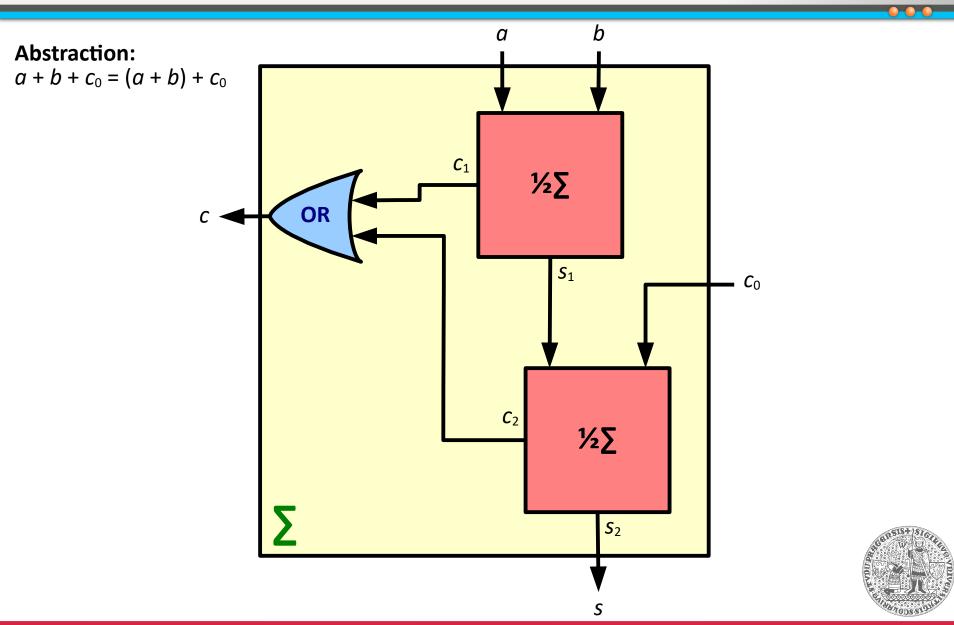


Logic blocks: binary adder (6)

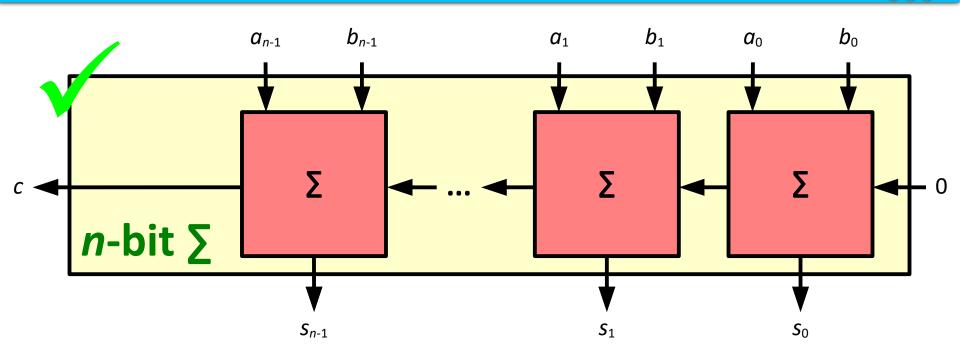


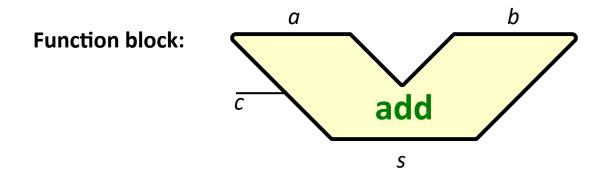


Logic blocks: binary adder (7)



Logic blocks: binary adder (8)







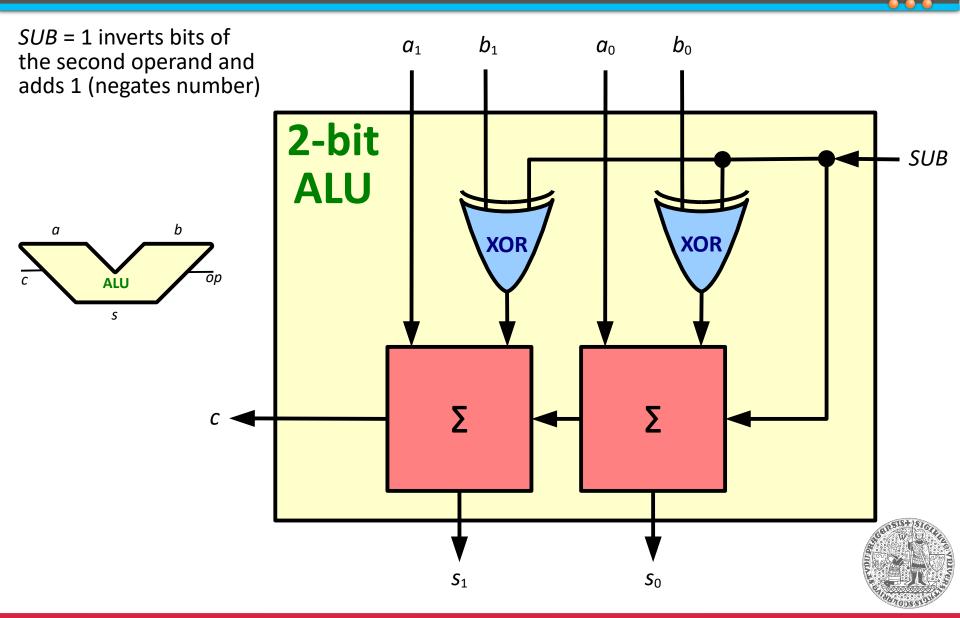
Logic block for subtraction

• Taking advantage of 2's complement

- Basic building block: adder
- Use XOR gate as a controlled inverter
- Example: 2-bit ALU supporting addition and subtraction
 - Data input: operand bits a₁a₀, operand bits b₁b₀
 - Control input: signal SUB to determine operation
 - $SUB = 0 \rightarrow addition$
 - $SUB = 1 \rightarrow \text{subtraction}$
 - **Output:** sum/difference bits *s*₁*s*₀, carry *c*



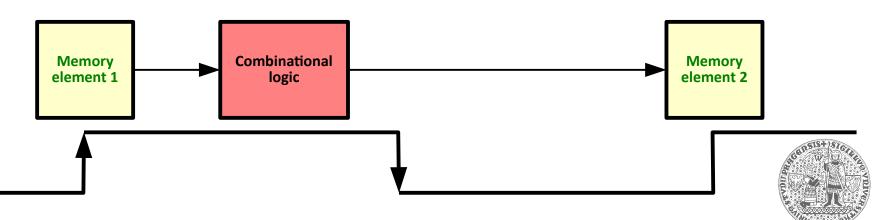
2-bit ALU for adding/subtracting



Sequential logic

• Combinational logic + memory elements

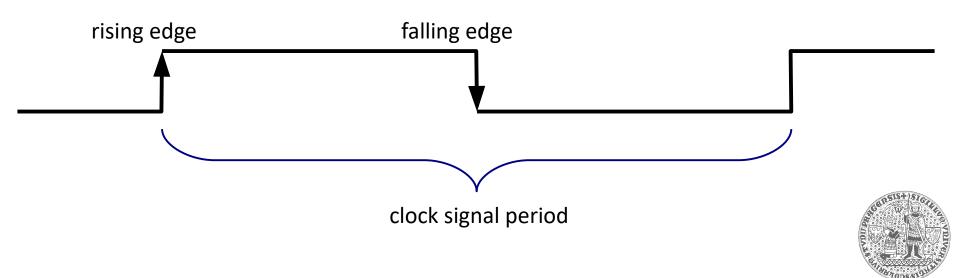
- Memory elements keep internal state
- Inputs and the contents of memory (internal state) determines outputs and next internal state
 - Synchronous vs. asynchronous sequential circuits
 - Determines how and when state changes
 - Need to ensure stable inputs (inputs don't change)



Synchronous sequential circuits

• Clock signal to synchronize state changes

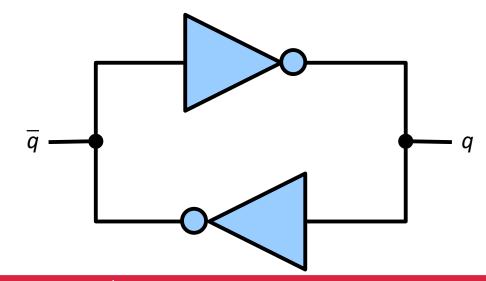
- Change state during one clock cycle
 - Inputs of combinational logic does not change while it is being read
 - Writing of values from outputs to memory elements happens with rising/falling edge of the clock signal



Memory elements

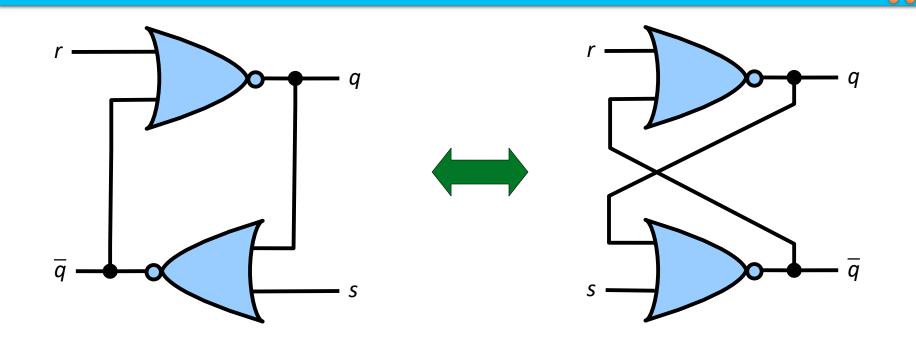
• Pair of inverters in a feedback loop

- Asynchronous circuit with two stable states
 - Allows "storing" 1 bit of information
 - Need to be able to control the state...
 - We need a gate that can pass the signal unchanged, but allows forcing an output value when required
- Basic building block for memory elements





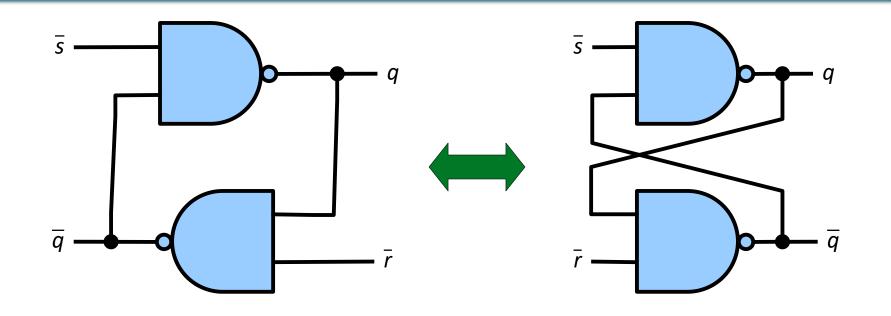
Set-Reset (R-S) latch, NOR-based



Inp	uts	Outputs		
r	S	\boldsymbol{q}_n	$\overline{oldsymbol{q}}_n$	
0	0	q _{n-1}	¬ <i>q</i> _{<i>n</i>-1}	
0	1	1	0	
1	0	0	1	
1	1	?	?	



Set-Reset (R-S) latch, NAND-based



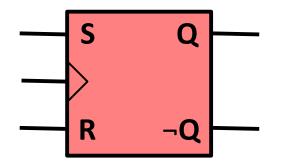
Inputs		Outputs		
r	Ś	\boldsymbol{q}_{n}	$\overline{\boldsymbol{q}}_{n}$	
0	0	?	?	
0	1	0	1	
1	0	1	0	
1	1	q _{n-1}	¬ q _{n-1}	



Other flip-flops

Derived from R-S

- Clocked R-S latch
 - Synchronous R-S latch variant
 - Reacts to **R** or **S** signals while the clock signal is high
- R-S master/slave (R-S flip-flop)
 - Two clocked R-S latches (in series) with complementary clock signal
 - Reacts to R or S signals only on rising/falling edge of the clock signal

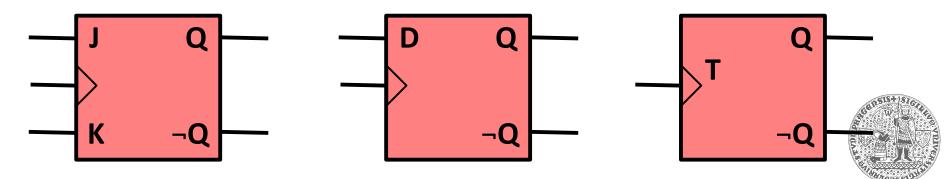




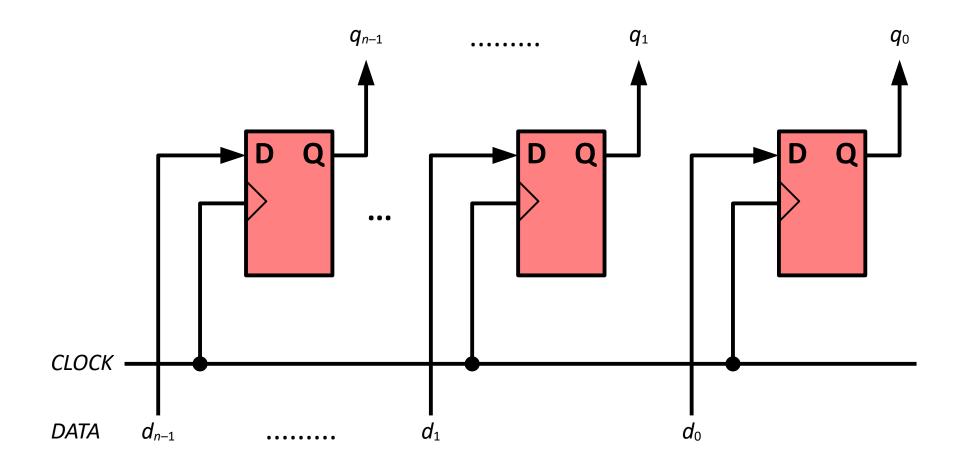
Other flip-flops (2)

Derived from R-S

- J-K master/slave (J-K flip-flop)
 - Extends R-S (J = S, K = R), inverts state when J = K = 1
- Clocked D latch, D flip-flop
 - Value determines by single input
- T flip-flop
 - Allows dividing clock signal frequency

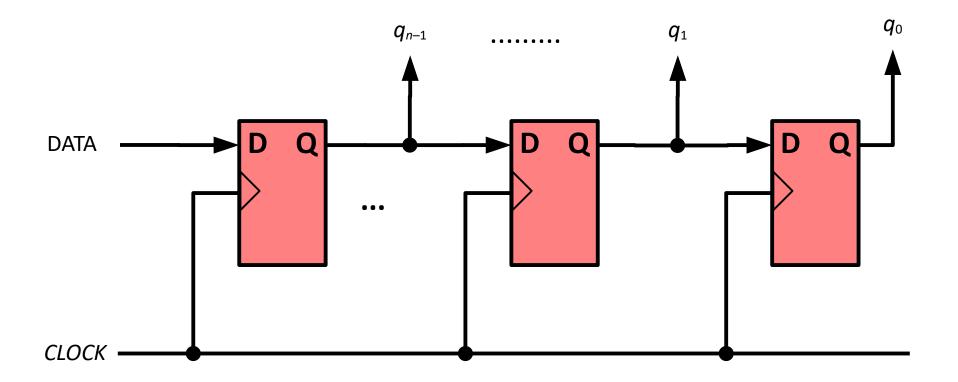


Data register made of flip-flops



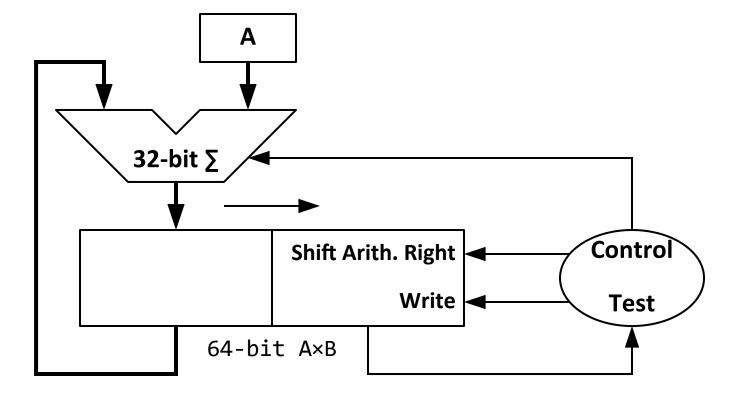


Shift register made of flip-flops





32-bit sequential multiplier





32-bit sequential divider

