

Computer (Literacy) Skills

Number representations and memory

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Number representations? What for?

Recall: computer works with binary numbers

- Groups of zeroes and ones
 - 8 bits (byte), 16 bits (2 bytes), ..., 64 bits (8 bytes), ...
 - Number of bits determines range of numbers
 - If a number does not fit \Rightarrow overflow!
- So far we have used **unsigned integer** numbers
 - Natural numbers with zero included

People usually need other kinds of numbers

- Positive, negative, fractions, ...
- Groups of zeroes and ones... again?

We can interpret bits in different ways

- Some interpretations more useful than others

Positive/negative integer numbers

Things to consider

- Distinguishing positive/negative numbers
- Difficulty working with the numbers (in HW)
 - Arithmetic, negation, detecting overflow
 - Extending/truncating fixed-size representation

Sign and magnitude representation

- Explicit sign bit (where to put it?)
- N bits $\Rightarrow \{ -(2^{N-1}), \dots, -\mathbf{0}, +\mathbf{0}, \dots, 2^{N-1} \}$

Biased representation

- Implicit sign bit (only for bias of $2^{N-1} - 1$)
- N bits with bias $B \Rightarrow \{ -B, \dots, \mathbf{0}, \dots, 2^N - 1 - B \}$

Positive/negative integer numbers

One's complement representation

- Implicit sign bit, symmetric range
- Negation is “flip all bits”
- N bits $\Rightarrow \{ -(2^{N-1}), \dots, -\mathbf{0}, +\mathbf{0}, \dots, 2^{N-1} \}$

Two's complement representation

- Implicit sign bit, asymmetric range
- Negation is “**flip all bits, then add 1**”
- N bits $\Rightarrow \{ -(2^{N-1}), \dots, \mathbf{0}, \dots, 2^{N-1} - 1 \}$
 - $b_{N-1} \times -(2^{N-1}) + b_{N-2} \times 2^{N-2} + \dots + b_1 \times 2^1 + b_0 \times 2^0$
- Modular arithmetic!
 - Subtract by adding in unsigned arithmetic

Extending/truncating numbers

General principle

- Ensure that interpreting the extended/truncated representation gives the same number
- Truncation can result in **loss of information!**

Unsigned integers

- Extend with zero bits to the left, truncation trivial

Sign and magnitude representation

- Strip sign, extend/truncate as unsigned, set sign

One's and two's complement representations

- Extend using the value of the highest bit, truncation trivial

Representing fractional numbers

Fixed point representation

- Decimal analogy: numbers 0 ... 99 divided by 10 allow representing 0.0, 0.1, 0.2, ..., 9.9
- In binary, the fractional part of a number is the sum of negative powers of 2
- Works also with two's complement

Example: 4.4 fixed-point representation

- Integral part: 0, 1, ..., 15
- Fractional part: $0 \times 0.0625, \dots, 15 \times 0.0625 = 0.9375$

$2^3 = 8$	$2^2 = 4$	$2^1 = 2$	$2^0 = 1$	$2^{-1} = 0.5$	$2^{-2} = 0.25$	$2^{-3} = 0.125$	$2^{-4} = 0.0625$
$b_7 = \text{MSB}$	b_6	b_5	b_4	b_3	b_2	b_1	$b_0 = \text{LSB}$



(Imaginary) fixed point

Decimal \leftrightarrow binary conversion

Fractional part

- Convert separately from integral part

Simple algorithm

- Multiply fractional part by 2
- Value before decimal point provides next fraction bit, starting with MSB
- Strip of the fractional part and repeat until zero, or ...
 - ... the pattern starts repeating
 - ... we have enough bits

$$0.678_{10} = ???_2$$

$$0.678 \times 2 = 1.356 \quad (1 = b_{-1})$$

$$0.356 \times 2 = 0.712 \quad (0 = b_{-2})$$

$$0.712 \times 2 = 1.424 \quad (1 = b_{-3})$$

$$0.424 \times 2 = 0.848 \quad (0 = b_{-4})$$

$$0.848 \times 2 = 1.696 \quad (1 = b_{-5})$$

$$0.696 \times 2 = 1.392 \quad (1 = b_{-6})$$

$$0.392 \times 2 = 0.784 \quad (0 = b_{-7})$$

$$0.784 \times 2 = 1.568 \quad (1 = b_{-8})$$

$$0.678_{10} \cong 0.10101101_2$$

Approximating real numbers

Floating point representation

- Decimal analogy: normalized scientific notation
 $D_0, D_1 D_2 \dots D_{P-1} \times 10^E$ (P valid digits, $1 \leq D_0 \leq 9$)

- Similarly in binary

$$B_0, B_1 B_2 \dots B_{P-1} \times 2^E \text{ (P valid bits, } 1 \leq B_0 \leq 1)$$

B_0 is always 1 in this form

In-memory representation

Sign	Exponent (with bias)	Significand (value of B_0 not stored \Rightarrow hidden 1)
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SP	DP
Bias = 127	Bias = 1023
P = 24	P = 53

- $(-1)^{\text{Sign}} \times \text{Significand} \times 2^{(\text{Exponent} - \text{Bias})}$
- Half (16-bit) / Single (32-bit) / Double (64-bit)

Real number → IEEE floating point

- 1. Convert to binary fractional number**
 - Integral and fractional part, ignore sign
- 2. Normalize the binary representation**
 - Move binary point to get $1.ssss \times 2^{\text{Exp}}$
- 3. Depending on the target FP representation**
 - Round significand to desired precision
 - Convert the exponent to biased representation
- 4. Set the sign bit to reflect the sign**
- 5. For in-memory representation**
 - Drop initial 1 from the significand (hidden 1)

Speaking of memory...

Programmer's perspective (logical view)

- 1-D array of N bytes numbered 0, ..., N-1
- Individual bytes can be read or written to
- A particular byte is identified by its index

Index => Address

- Numbers occupy multiple bytes in memory
- Address of something = address of first byte

Memory is “visible” to CPU

- CPU sends address to memory controller, requesting bytes to be read or written
- Memory controller uses parts of the address to determine which part of memory to access

How to store multi-byte numbers?

Memory = array of bytes

- Chop up number into sequence of N bytes
 - B_{N-1} (Most Signif. Byte), ..., B_1 , B_0 (Least Signif. Byte)
- Store N bytes at consecutive addresses in memory
 - Addresses A , $A+1$, $A+2$, ..., $A+(N-1)$ for N-byte number
- CPU does this when storing/loading contents of its registers to/from memory at a given address

The order of bytes matters!

- Similar to sending bits over serial line.

Big Endian = MSB first

A	A+1	...	A+(N-1)
B_{N-1}	B_{N-2}	...	B_0

Little Endian = LSB first

A	A+1	...	A+(N-1)
B_0	B_1	...	B_{N-1}

Understanding a memory dump

Lists contents of memory

- Or any other address space, e.g., storage device (hard disk, solid state drive), or a file
- Contents of starting address and fixed number of consecutive addresses (usually all in hexadecimal)

We must know the interpretation!

- What is stored at 0x03194A7B? Or at 0x03194A84?

Address	Byte at (Address + 0), (Address + 1), ..., (Address + 7)							
...								
03194A78	AF	BC	39	F6	D0	24	91	34
03194A80	81	C9	A3	7C	00	80	B7	C2
03194A88	E2	6C	71	EA	59	FE	F5	49
...								