

Symbolic Execution, Dynamic Analysis

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Symbolic execution



Key concepts

- Symbolic values used for program variables
 - $v: x+2, x: i0+i1-3, y: 2*i1$
- Program inputs: variable names
- Other variables: functions over symbolic inputs
- Path condition (PaC)
 - Set of constraints over symbolic input values that hold in the current program state

Example program

input: x, y

1: $u = x - y;$

2: $\text{if } (x > 1)$

3: $u = u + x;$

4: $\text{if } (y > x)$

5: $u = y - x;$

6: $\text{assert } (u \geq 0);$

Symbolic execution and program verification

- Symbolic program state
 - Symbolic values of program variables
 - Path condition (PaC)
 - Program counter (PC)
- Symbolic state = a set of concrete states
- Symbolic execution tree = state space
 - Tree of symbolic program states
 - Transitions labeled with the PC

Symbolic execution and program verification

- Path condition updated at each branching point in the program code
 - Different constraints added for each branch
 - Example: `if-else` with a boolean condition C
 - Formula C added for the `if` branch
 - Formula $\text{not } C$ added for the `else` branch
- State space traversal
 - Satisfiability of PaC checked in each symbolic state
 - $PaC == \text{false}$ \rightarrow symbolic state not reachable
 - Verification tool backtracks and explores different branches

Symbolic execution: possible applications



Do you have some ideas ?

Symbolic execution: possible applications

- Automatically generating test inputs
 - From path conditions in symbolic program states
- Find inputs that trigger a specific error
- Systematic testing of open systems
 - Examples: isolated procedures, components
 - Programs with unspecified concrete inputs
- Checking programs with inputs from unbounded domains (integers, floats, strings)

Symbolic execution: limitations



What are the limitations ?

Symbolic execution: limitations

- Handling loops with many iterations
- Stateless exploration (no state matching)
- Undecidable and complex path conditions
- State explosion (too many paths)
- Concurrent accesses from multiple threads

Loops with many iterations

```
x = input();  
i = 1000;  
while (true) {  
    if (x > i) ...  
    i--;  
    ...  
}
```

Loops: practical approach

- Unrolling loops to a specific depth
 - Limited number of loop iterations explored
- Exploring data structures up to a given bounded size

Concolic execution



Concolic execution

concrete + symbolic = concolic

- How it works
 - Performs concrete execution on random inputs
 - Tracks symbolic values of program variables
 - Gathers constraints forming a path condition along the single executed path

Concolic execution: applications

- Dynamic test generation
 - Path condition for the single explored path defines corresponding test inputs
 - Negating constraints (clauses) for branching points
 - Find test inputs that drive program execution along different paths (control-flow)

KLEE: Symbolic Virtual Machine

- Symbolic execution tool for system code
 - Used to detected many real bugs in Linux/Unix core system utilities (`ls`, `chmod`, ...)
 - Models interaction with complex environment (files, networking, unix syscalls)
 - Highly optimized (performance, scalability)
- Built upon the LLVM compiler infrastructure
- Web: <http://klee.github.io/>
- Further information (recommended)
 - <http://llvm.org/pubs/2008-12-OSDI-KLEE.pdf>

PEX: White-box unit testing for .NET

- Dynamic test generation (concolic execution)
- Generates unit tests with high code coverage
- Availability
 - Visual Studio 2010 Power Tools, command-line
- Web sites
 - <https://www.microsoft.com/en-us/research/project/pex-and-moles-isolation-and-white-box-unit-testing-for-net/>
- Live demo: Code Digger
 - Visual Studio 2012 extension based on Pex
- IntelliTest extension for Visual Studio 2015
 - <https://docs.microsoft.com/en-us/visualstudio/releases/notes/vs2015-rtm-vs#intellitest>

SAGE: Scalable Automated Guided Execution

- Automated whitebox fuzz testing for security
- Systematic dynamic generation of unit tests
- How it works
 - concolic execution + solving negated conditions to infer new test inputs
- Main author: Patrice Godefroid
 - <https://patricegodefroid.github.io/>
- Further information (selected papers)
 - https://patricegodefroid.github.io/public_psfiles/ndss2008.pdf
 - https://patricegodefroid.github.io/public_psfiles/icse2013.pdf

Other tools

- <https://playground.diffblue.com/>
 - Automatically generating tests for Java code
 - Registration needed (or sign-in via FB/Google)

Dynamic analysis



Dynamic analysis

- Goal: analyze behavior of the program based on concrete execution of a single path
- Input: binary executable

Q: How can we get the data about one path?

Collecting information about single path

- Instrumentation
 - Target: binary executable, source code
- Runtime monitoring
 - manual inspection of huge log files
- Custom libraries
- Events
 - field accesses on shared heap objects
 - locking (acquisition, release, attempts)
 - procedure calls (e.g., user-defined list)

Benefits

- Precision
 - Complete information about program state
 - Recording only events that really happen
- Tool support
 - Errors: deadlocks, race conditions, atomicity
 - Languages: Java, C/C++, C# ?

Limitations

- Coverage
 - Single execution path
 - Few related paths
- Overhead
 - Compared with plain concrete execution
 - Range: 50 % - 1000 % (!)
 - Possible remedy: sampling

Selected tools

- Pin
 - Runtime binary instrumentation platform for Linux (32-bit x86, 64-bit x86, ARM)
 - Custom tools written in C/C++ using rich Pin API
 - Important features
 - efficient dynamic compilation (JIT)
 - process attaching, transparency
- Valgrind
 - Heavyweight dynamic binary instrumentation framework again for Linux (x86, PPC)
 - Tools: memory checker, thread checkers, some profilers
- RoadRunner

Applications

- Detecting bugs of all kinds
- Concolic execution
 - Adding new symbolic constraints into PaC
- Discovering likely invariants
- Predicting race conditions

Predicting data race conditions

- Algorithm

- 1) Run dynamic analysis tool to record events about one particular execution trace
- 2) Check the given trace for data race conditions
- 3) If we find some errors, then stop immediately
- 4) Generate feasible interleavings of events from the given single trace
- 5) Check each generated interleaving for data races
- 6) Report all detected possible races to the user

Predicting data race conditions



Q: How can we generate feasible interleavings ?

Generating feasible interleavings

- All possible interleavings of events from different threads
- Use the happens-before order between synchronization events
 - Conflicting field accesses not ordered → interleave

Discovering likely invariants

- Algorithm
 - Run the dynamic analysis tool several times (on selected inputs, test suite) to get a set of traces
 - Find properties over variables and data structures that hold for all/most traces in the set
 - Drop all inferred properties that do not satisfy additional tests (e.g., statistical relevance)
 - What remains are the likely invariants

Q: Looks good but there is a small catch

Discovering likely invariants

- Limitations
 - Precision depends on the test suite quality (inputs)
 - Cannot guarantee soundness and completeness
- Benefits
 - It is actually useful: checking implicit assumptions about program behavior, rediscovering formal specifications, documentation, etc
- Tool support: Daikon
 - Predefined templates instantiated with variables
 - <http://plse.cs.washington.edu/daikon/>

Further reading

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